

A proposal for parallel implicit time evolution

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joint work with
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Heat equation:

$$\frac{\partial u}{\partial t} - \nabla^2 u = f \quad \text{in } \Omega \times (0, T], \quad \Omega \subset \mathbb{R}^2 \text{ or } \mathbb{R}^3$$

$$u(\mathbf{x}, 0) = u_0(\mathbf{x}), \quad u \text{ given on } \partial\Omega$$

Finite elements in space (\mathbf{x}), θ time stepping gives

$$M \frac{\mathbf{u}_k - \mathbf{u}_{k-1}}{\tau} + K \left(\theta \mathbf{u}_k + (1 - \theta) \mathbf{u}_{k-1} \right) = \mathbf{f}_k$$

$M \in \mathbb{R}^{n \times n}$: SPD mass matrix (identity operator, but same sparsity as K)

$K \in \mathbb{R}^{n \times n}$: SPD discrete (negative) Laplacian (stiffness matrix)

Rearranging:

$$\left(M + \tau \theta K \right) \mathbf{u}_k = \left(M - \tau (1 - \theta) K \right) \mathbf{u}_{k-1} + \tau \mathbf{f}_k,$$

$$k = 1, 2, \dots, N$$

$$N\tau = T$$

Recall for unconditional stability: $\frac{1}{2} \leq \theta \leq 1$

$\theta = 1$: backwards Euler, $\theta = \frac{1}{2}$: Crank-Nicolson

else need $\tau = \mathcal{O}(h^2)$: very small time steps for explicit method

$$\left(M + \tau \theta K \right) \mathbf{u}_k = \left(M - \tau (1 - \theta) K \right) \mathbf{u}_{k-1} + \tau \mathbf{f}_k,$$

$$k = 1, 2, \dots, N$$

Standard solution method:

Solve the N separate $n \times n$ linear systems **sequentially** for $k = 1, 2, \dots, N$ e.g. by algebraic multigrid (we use HSL_MI20)

$\Rightarrow r = 5$ V-cycles for solution of each linear system to a relative residual tolerance of 10^{-6}

Hence if we (quite reasonably) regard 1 V-cycle as the main unit of work

$\Rightarrow Nr$ V-cycles sequentially for the overall solution

Alternative proposal for parallel computation:

Write all timesteps at one go (all-at-once method):

$$\mathcal{A} \begin{bmatrix} \mathbf{u}_1 \\ \mathbf{u}_2 \\ \vdots \\ \mathbf{u}_N \end{bmatrix} = r.h.s$$

where \mathcal{A} is the matrix

$$\begin{bmatrix} M+\tau\theta K & 0 & 0 & 0 \\ -M+\tau(1-\theta)K & M+\tau\theta K & 0 & 0 \\ 0 & \ddots & \ddots & 0 \\ 0 & 0 & -M+\tau(1-\theta)K & M+\tau\theta K \end{bmatrix}$$

and $r.h.s. = [M-\tau(1-\theta)K \mathbf{u}_0 + \tau \mathbf{f}_1, \tau \mathbf{f}_2, \dots, \tau \mathbf{f}_N]^T$

$$\mathcal{A} = \begin{bmatrix} M + \tau\theta K & 0 & 0 & 0 \\ -M + \tau(1 - \theta)K & M + \tau\theta K & 0 & 0 \\ 0 & \ddots & \ddots & 0 \\ 0 & 0 & -M + \tau(1 - \theta)K & M + \tau\theta K \end{bmatrix}$$

$$\mathcal{A} \in \mathbb{R}^{L \times L}, \quad L = Nn$$

We propose to solve this huge linear system (for the solution at all time steps) by GMRES (or BICGSTAB) with block diagonal preconditioner

$$\mathcal{P} = \begin{bmatrix} (M + \tau\theta K)_{MG} & 0 & 0 & 0 \\ 0 & (M + \tau\theta K)_{MG} & 0 & 0 \\ 0 & \ddots & \ddots & 0 \\ 0 & 0 & 0 & (M + \tau\theta K)_{MG} \end{bmatrix}$$

where $(M + \tau\theta K)_{MG}$ is one AMG V-cycle exactly as above

Theory: If we used

$$\mathcal{P}_{\text{exact}} = \begin{bmatrix} (M+\tau\theta K) & 0 & 0 & 0 \\ 0 & (M+\tau\theta K) & 0 & 0 \\ 0 & \ddots & \ddots & 0 \\ 0 & 0 & 0 & (M+\tau\theta K) \end{bmatrix}$$

as preconditioner (no AMG approximation) then we would have

$$\mathcal{P}_{\text{exact}}^{-1} \mathcal{A} = \begin{bmatrix} I & 0 & 0 & 0 \\ J & I & 0 & 0 \\ 0 & \ddots & \ddots & 0 \\ 0 & 0 & J & I \end{bmatrix},$$

$$J = (M+\tau\theta K)^{-1} (-M+\tau(1-\theta)K)$$

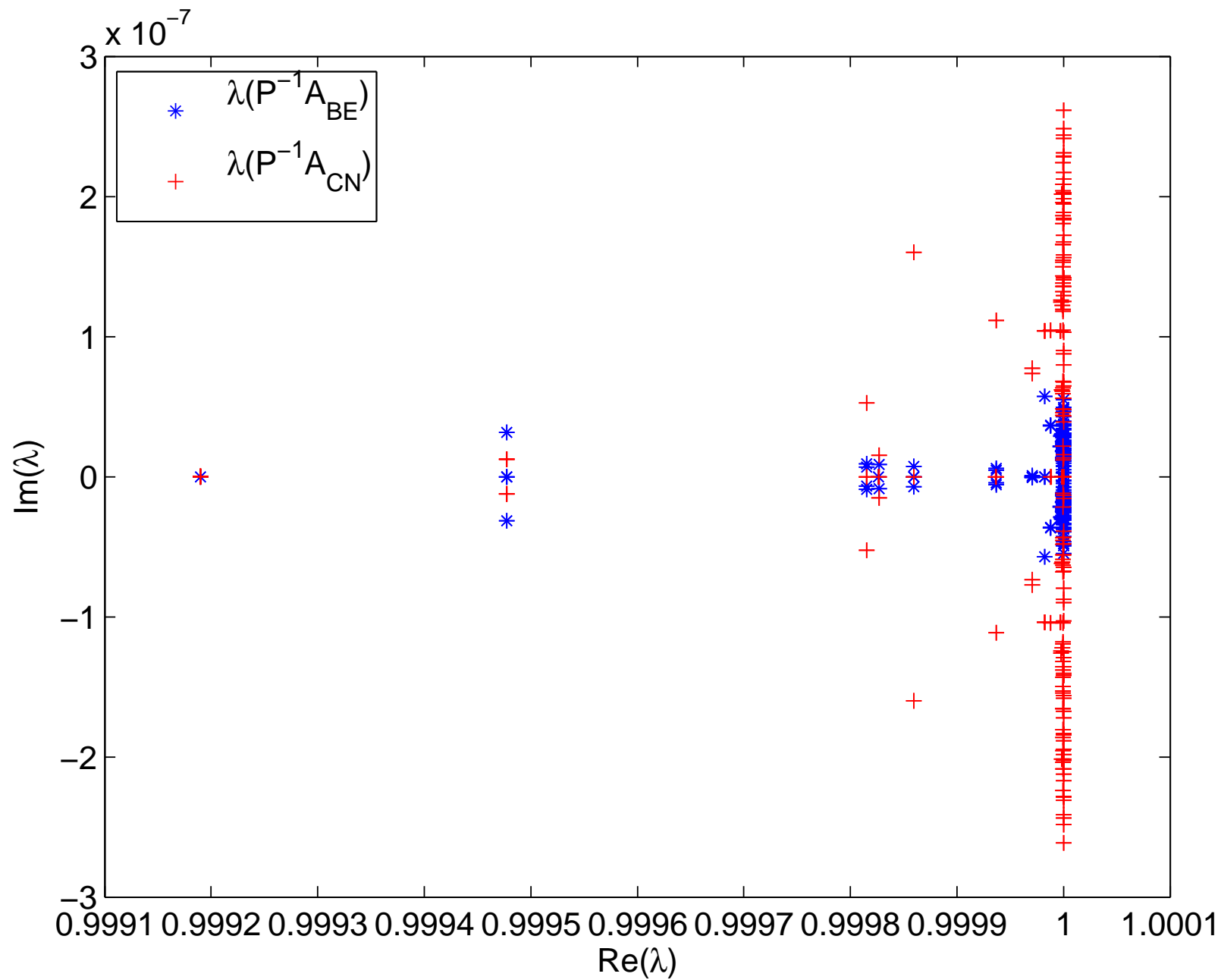
For

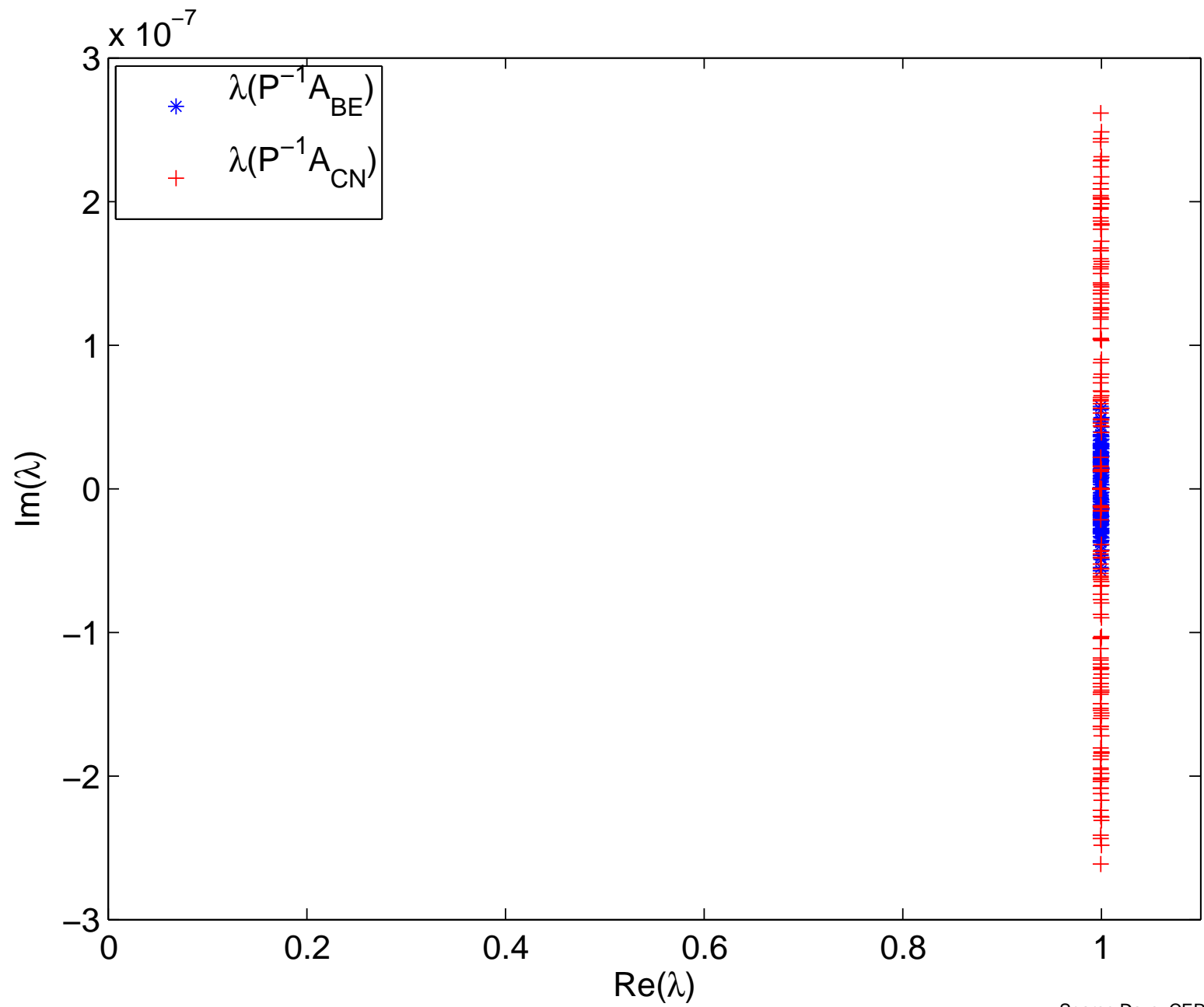
$$\mathcal{P}_{\text{exact}}^{-1} \mathcal{A} = \begin{bmatrix} I & 0 & 0 & 0 \\ J & I & 0 & 0 \\ 0 & \ddots & \ddots & 0 \\ 0 & 0 & J & I \end{bmatrix},$$

the minimum polynomial is $(1 - s)^N$, so GMRES would terminate (in exact arithmetic) in N iterations

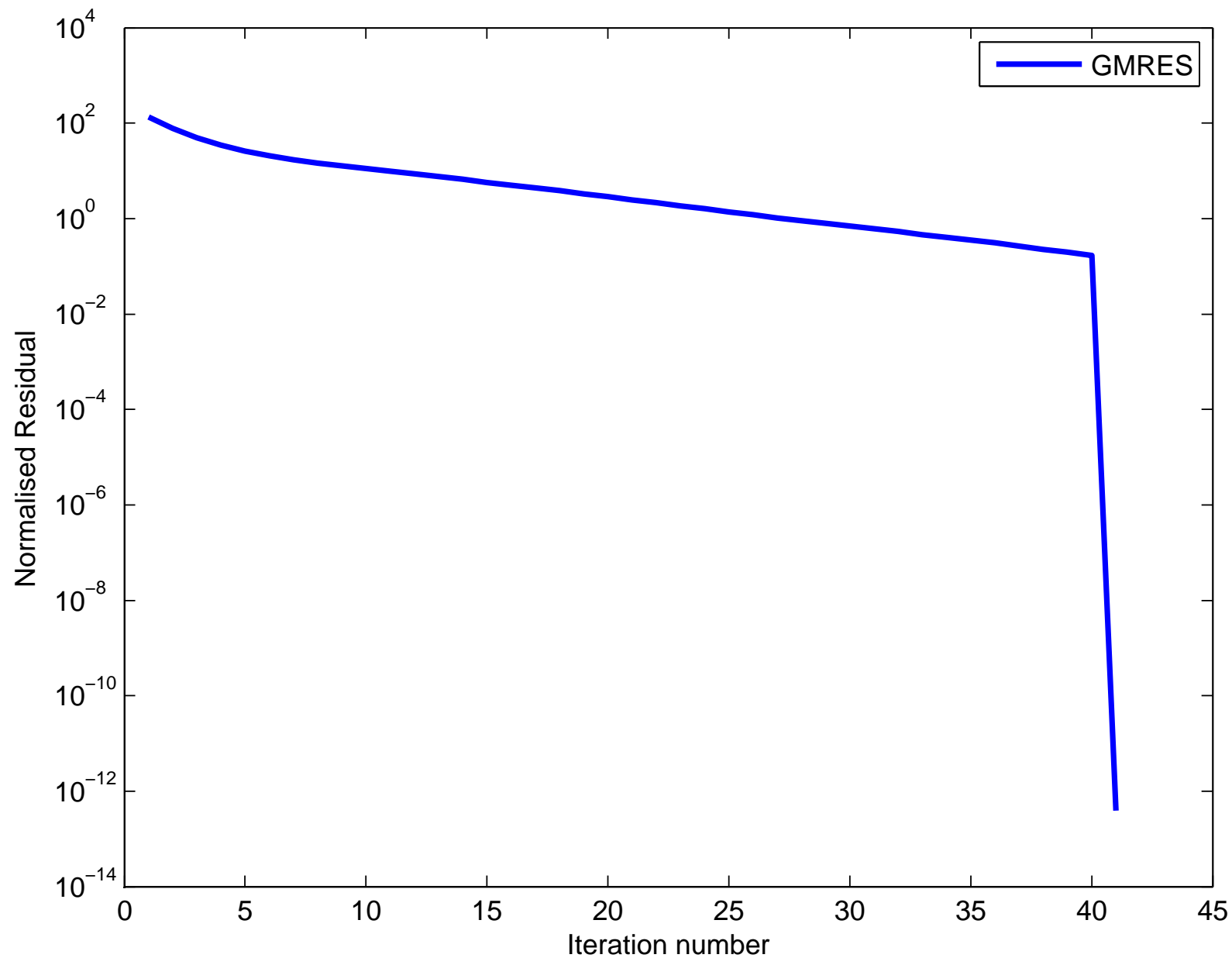
We observe that $(M + \tau\theta K)_{MG}$ is spectrally so close to $(M + \tau\theta K)$ that convergence to a tolerance much less than the discretization error is achieved in N iterations also with \mathcal{P} as preconditioner.

For N=5:





For N=40:

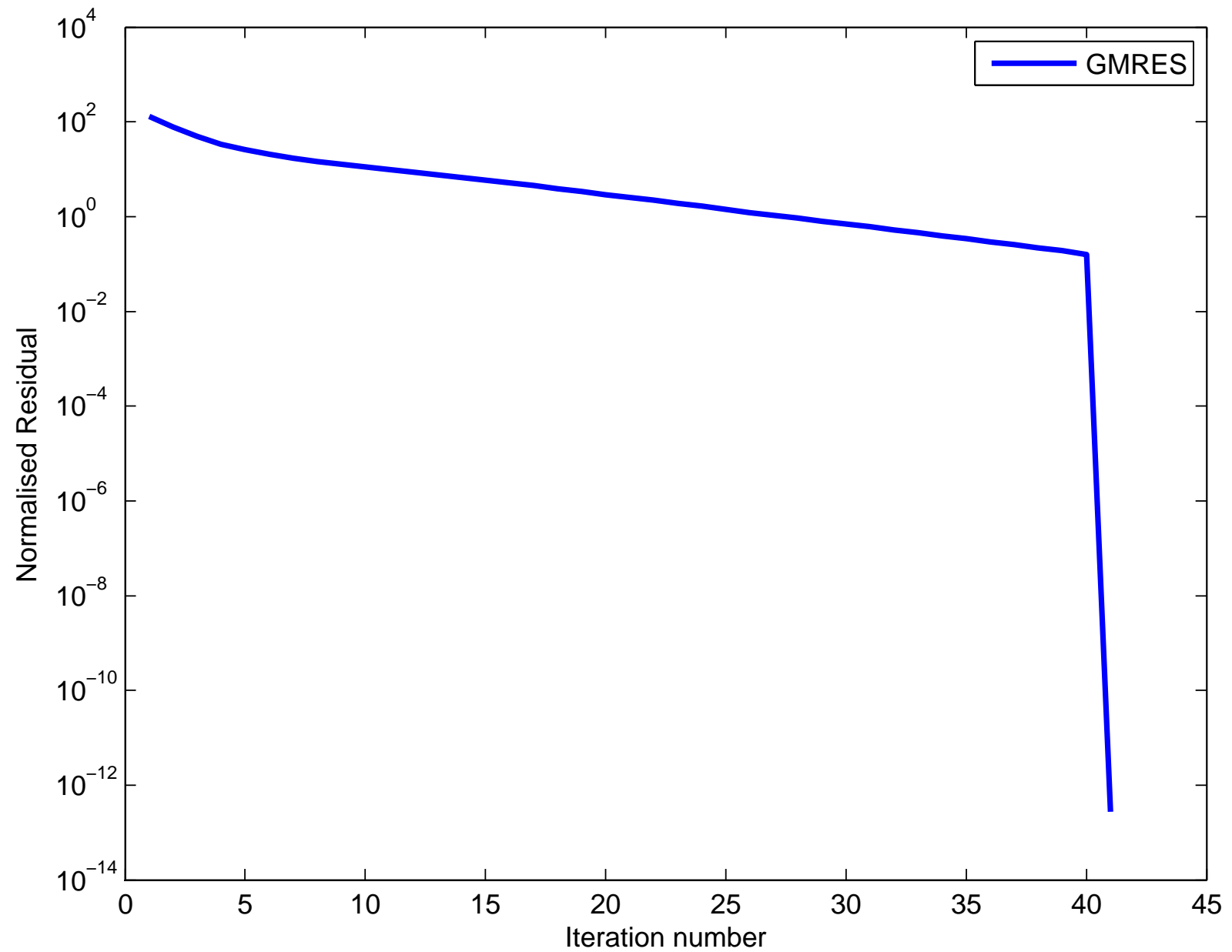


Thus: N V-cycles for each of N GMRES iterations—hence N^2 ($> Nr$) overall.

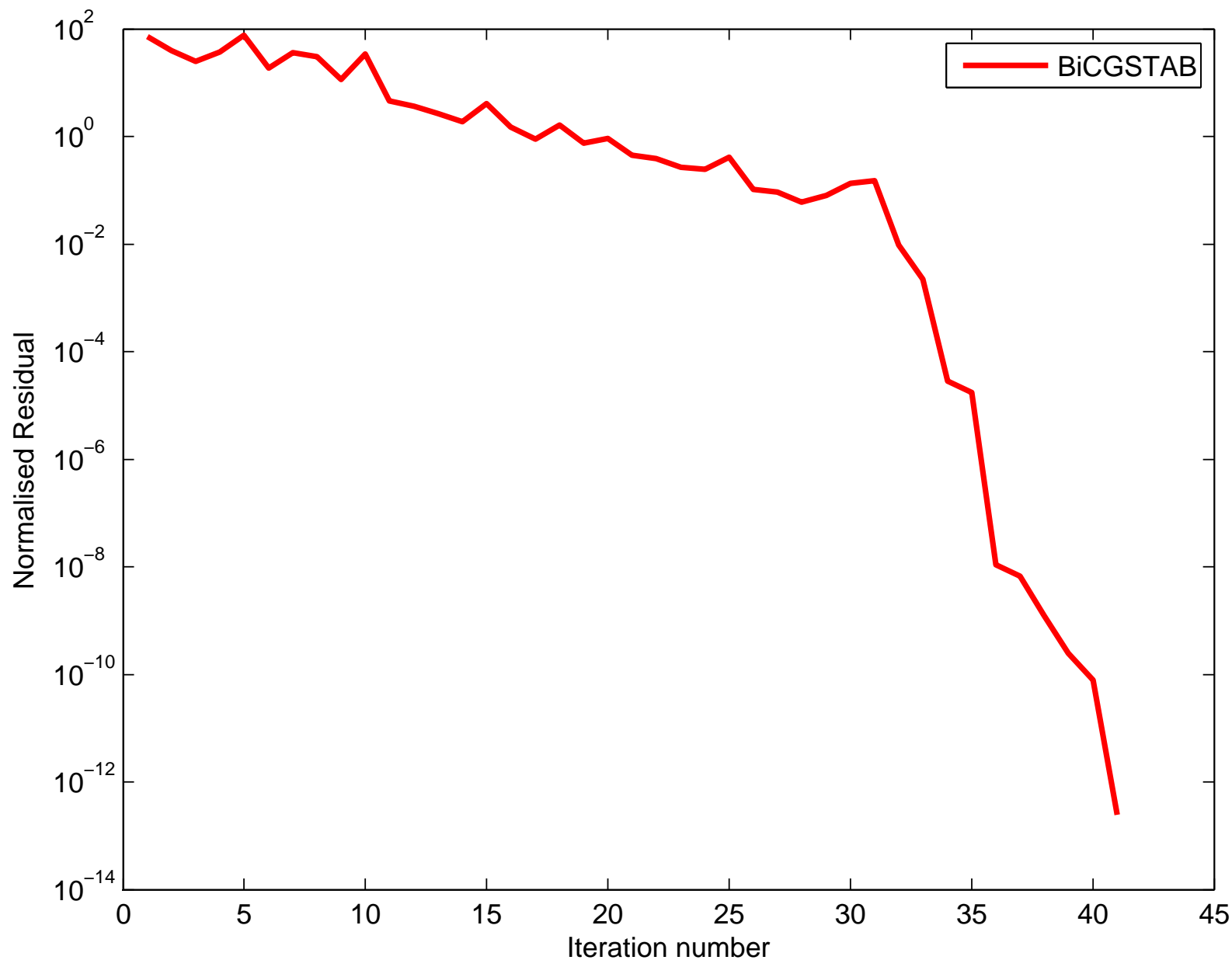
BUT with N processors, solution with \mathcal{P} is (embarrassingly) parallel—block diagonal \Rightarrow independent computation.

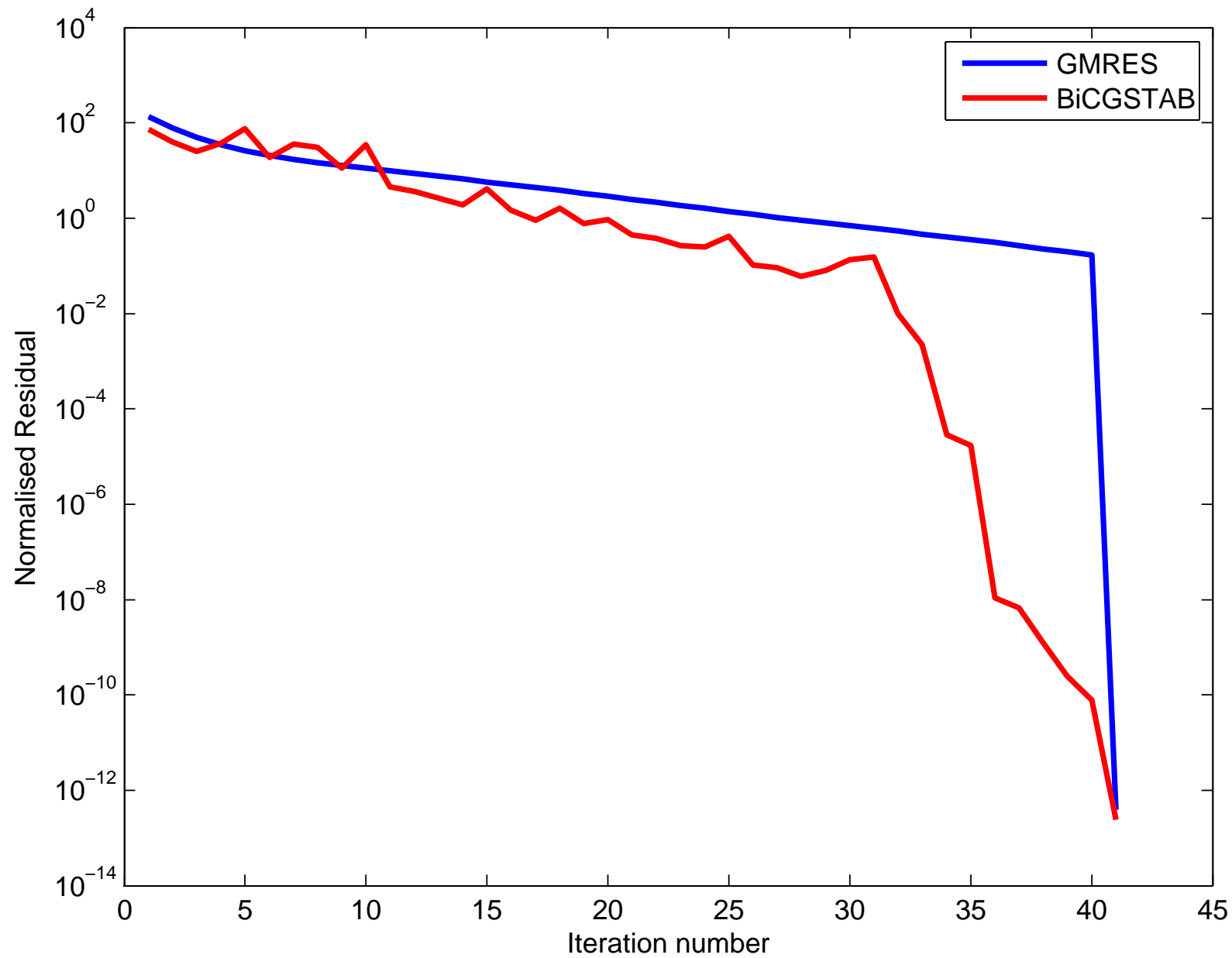
Thus parallel effort is $N < Nr$ (= sequential effort).

For another problem:

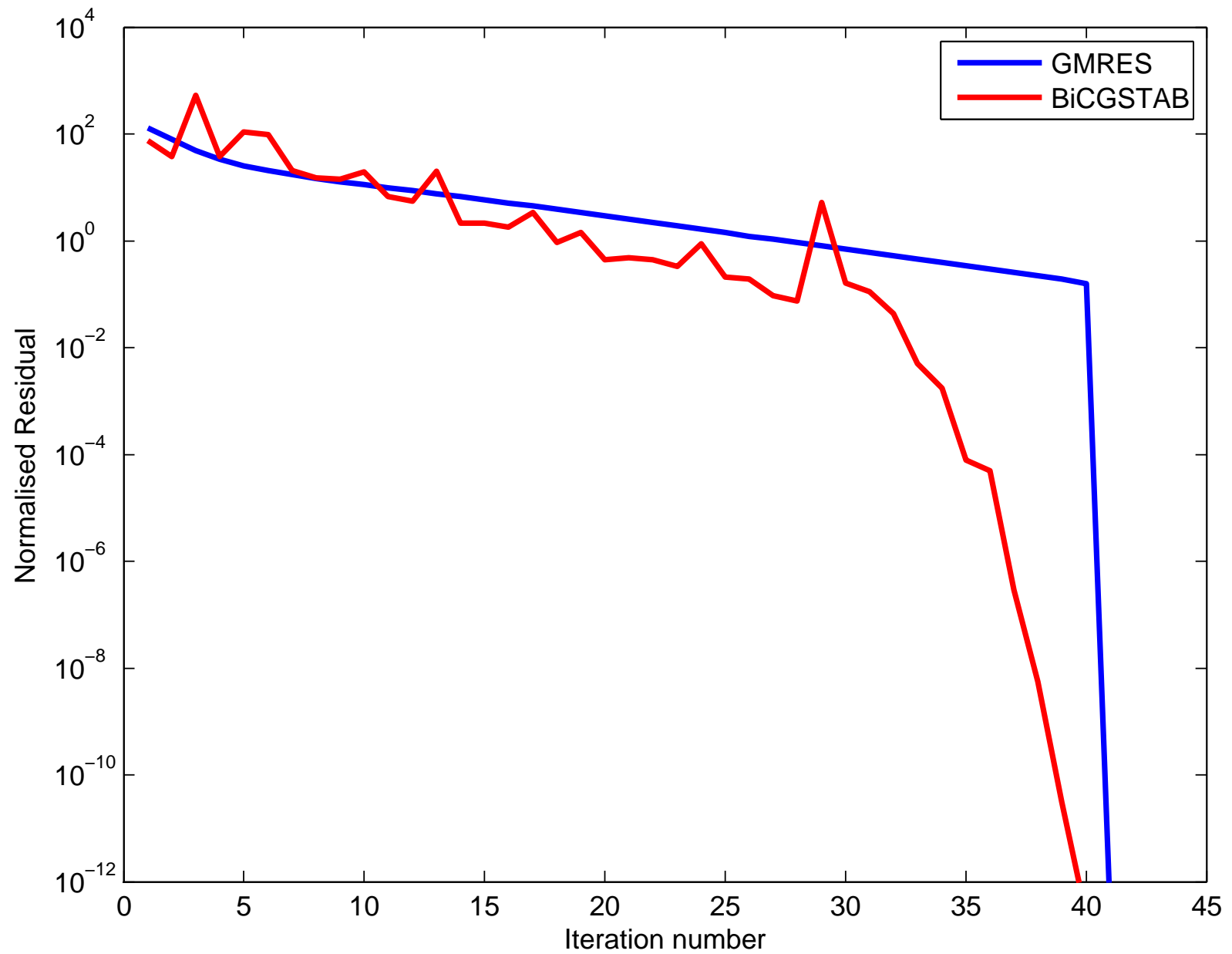


Large N means increasing work (of Arnoldi orthogonalisation) with GMRES : use BICGSTAB :





For another problem:



Backwards Euler

h	τ	N	DoF	GMRES	BiCGStab
2^{-3}	2^{-3}	40	3240	40	38
2^{-4}	2^{-4}	80	23120	80	78
2^{-5}	2^{-5}	160	174240	160	157
2^{-3}	2^{-3}	40	3240	40	38
2^{-4}	2^{-4}	40	11560	40	40
2^{-5}	2^{-5}	40	43560	40	43
2^{-6}	2^{-6}	40	169000	43	44
2^{-7}	2^{-7}	40	665640	45	45
2^{-8}	2^{-8}	40	2641960	46	45

Crank-Nicholson

h	τ	N	DoF	GMRES	BiCGStab
2^{-3}	2^{-5}	32	2592	33	35
2^{-4}	2^{-6}	64	18496	66	68
2^{-5}	2^{-7}	128	139392	132	138
2^{-3}	2^{-5}	32	2592	34	35
2^{-4}	2^{-6}	32	9248	35	34
2^{-5}	2^{-7}	32	34848	37	35
2^{-6}	2^{-8}	32	135200	39	36
2^{-7}	2^{-9}	32	532512	40	38
2^{-8}	2^{-10}	32	2113568	38	39

Summary

For a simple linear PDE problem our proposal should achieve

N work on N processors (N^2 work on 1 processor)

compared to

Nr work for the standard sequential algorithm

$\frac{1}{2}Nr$ work (?) for Parareal (?)

Acknowledgement

This work is partially supported by Award No. KUK-C1-013-04 made by King Abdullah University of Science and Technology (KAUST)