

A recursive model-based trust-region method for derivative-free bound-constrained optimization.

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Outline

- 1 Introduction
- 2 Interpolation models and poisedness
- 3 Geometry control in DFO trust region methods
- 4 Extension to bounds
- 5 Numerical experiments
- 6 Conclusions and Perspectives

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Why using derivative-free optimization?

Some reasons to apply Derivative-Free Optimization (DFO):

- Derivatives are unavailable
- Function evaluations are costly and/or noisy - Accurate approximation of derivatives by finite differences is prohibitive
- Source code not available or owned by a company - Automatic differentiation impossible to apply
- Growing sophistication of computer hardware and mathematical algorithms and software (opens new possibilities for optimization)

Applications:

- Tuning of algorithmic parameters
- Medical image registration
- Engineering design optimization, ...

Problem formulation

We consider the **bound-constrained minimization problem**

$$\min_{x \in \mathbb{R}^n} f(x) \quad \text{s.t.} \quad xl(i) \leq x(i) \leq xu(i), i = 1, \dots, n$$

where the first derivatives of the objective function are assumed to exist and be Lipschitz continuous, although explicit evaluation of these derivatives is assumed to be impossible.

We consider a **model-based trust-region algorithm** for computing local solutions of the minimization problem.

The method iteratively uses a **local interpolation model** of the objective function $f(x)$ to define a descent step, and adaptively adjusts the region in which this model is trusted.

Bibliography on developments in model-based DFO

Numerical optimization using local models:

- Powell, "A direct search optimization method that models the objective function by quadratic interpolation", 1994
- Conn, Scheinberg, and Toint, "On the convergence of derivative-free methods for unconstrained optimization", 1996
- Powell, "The NEWUOA software for unconstrained optimization without derivatives", 2004
- Conn, Scheinberg, and Vicente, "Introduction in Derivative Free Optimization", 2008
- Fasano, Nocedal, and Morales, "On the geometry phase in model-based algorithms for derivative-free optimization", 2009
- Scheinberg and Toint, "Self-correcting geometry in model-based algorithms for derivative-free unconstrained optimization", 2009

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Polynomial interpolation

Consider \mathcal{P}_n^d , the space of polynomials of degree $\leq d$ in \mathbb{R}^n .

A **polynomial basis** $\phi = \{\phi_1(x), \phi_2(x), \dots, \phi_p(x)\}$ of \mathcal{P}_n^d is a set of p polynomials of degree $\leq d$ that span \mathcal{P}_n^d .

For any basis ϕ , any polynomial $m(x) \in \mathcal{P}_n^d$ can be written as

$$m(x) = \sum_{j=1}^p \alpha_j \phi_j(x),$$

where α_j are real coefficients.

Polynomial interpolation

Given a **sample set** $Y = \{y^1, y^2, \dots, y^p\} \subset \mathbb{R}^n$ and a polynomial $m(x)$ of degree d in \mathbb{R}^n that interpolates $f(x)$ at the points Y , the coefficients $\alpha_1, \dots, \alpha_p$ **can be determined** by solving the linear system

$$M(\phi, Y)\alpha_\phi = f(Y),$$

where

$$M(\phi, Y) = \begin{bmatrix} \phi_1(y^1) & \phi_2(y^1) & \cdots & \phi_p(y^1) \\ \phi_1(y^2) & \phi_2(y^2) & \cdots & \phi_p(y^2) \\ \vdots & \vdots & & \vdots \\ \phi_1(y^p) & \phi_2(y^p) & \cdots & \phi_p(y^p) \end{bmatrix}, \quad f(Y) = \begin{bmatrix} f(y^1) \\ f(y^2) \\ \vdots \\ f(y^p) \end{bmatrix}.$$

Poisedness

If the coefficient matrix $M(\phi, Y)$ of the system is **nonsingular**, the set of points Y is called **poised** for polynomial interpolation in \mathbb{R}^n , otherwise the set Y is called non-poised.

As poisedness alone doesn't define the distance from singularity, there exists a measure of **well-poisedness**.

The most commonly used measure of **well-poisedness** in the polynomial interpolation literature is based on **Lagrange polynomials** [Powell, 1994].

Lagrange polynomials

If the sample set Y is **poised**, the basis of Lagrange polynomials **exists** and is **uniquely defined** (and vice versa).

The unique polynomial $m(x)$ that **interpolates** $f(x)$ on Y using the basis of Lagrange polynomials for Y can be expressed as

$$m(x) = \sum_{i=1}^p f(y^i) \ell_i(x),$$

where

$$\ell_j(y^i) = \begin{cases} 1 & \text{if } i = j, \\ 0 & \text{if } i \neq j. \end{cases}$$

is the basis of **Lagrange polynomials**.

Lagrange polynomials - Illustration

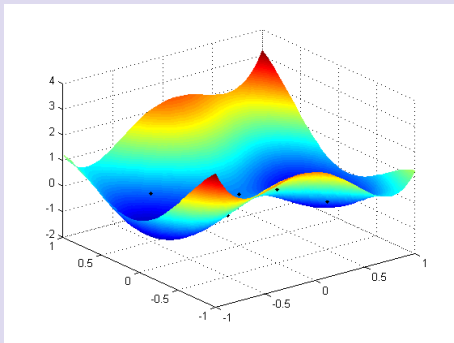


Figure: Six-hump camel back function with sample points

Lagrange polynomials - Illustration

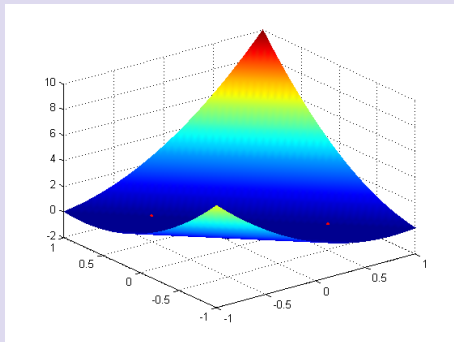


Figure: First Lagrange polynomial

Lagrange polynomials - Illustration

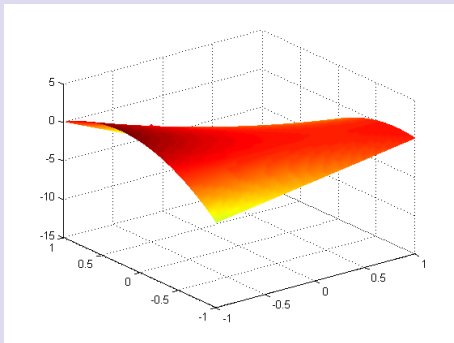


Figure: Second Lagrange polynomial

Lagrange polynomials - Illustration

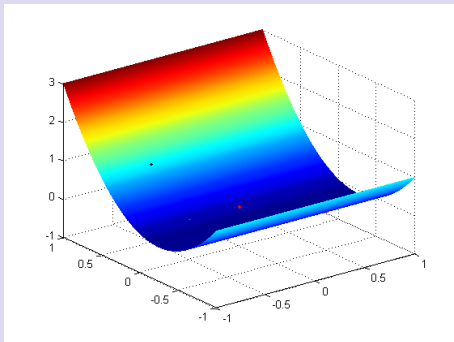


Figure: Third Lagrange polynomial

Lagrange polynomials - Illustration

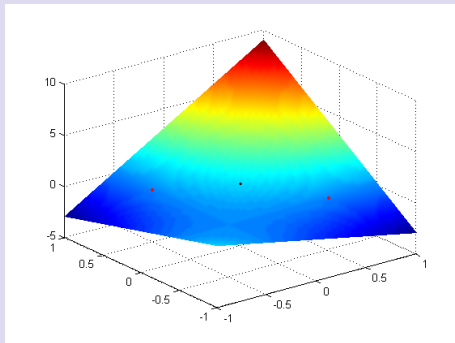


Figure: Fourth Lagrange polynomial

Lagrange polynomials - Illustration

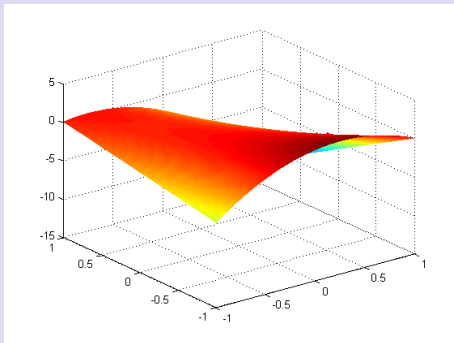


Figure: Fifth Lagrange polynomial

Lagrange polynomials - Illustration

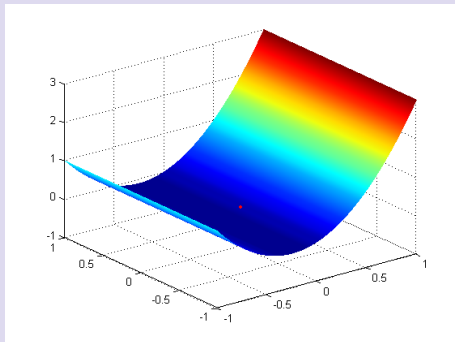


Figure: Sixth Lagrange polynomial

Lagrange polynomials - Illustration

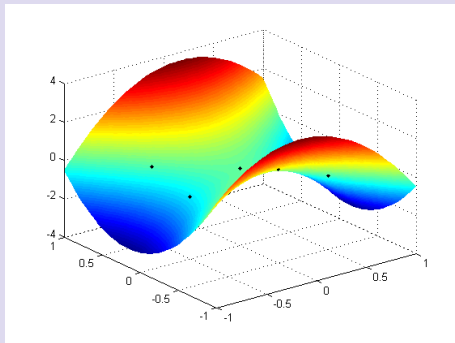


Figure: The resulting interpolation polynomial

Well poisedness

Very useful feature of Lagrange polynomials:

The upper bound on their absolute value in a region \mathcal{B} is a **classical measure of well-poisedness** of the interpolation set Y in the ball \mathcal{B} .

A poised set Y is said to be **Λ -poised in \mathcal{B}** if one has that

$$\max_{1 \leq i \leq p} \max_{x \in \mathcal{B}} |\ell_i(x)| \leq \Lambda.$$

The **smaller Λ** , the better the **quality of the geometry** of the interpolation set.

Error bounds on model value and model gradient value

Given a ball $\mathcal{B}(x, \Delta)$, a poised interpolation set $Y \in \mathcal{B}(x, \Delta)$ and its associated basis of Lagrange polynomials $\ell_i(x), i = 0, \dots, p$, there exists constants $\kappa_{ef} > 0$ and $\kappa_{eg} > 0$ such that, for any interpolation polynomial $m(x)$ of degree one or higher and any given point $y \in \mathcal{B}(x, \Delta)$,

$$\|f(x) - m(x)\| \leq \kappa_{ef} \sum_{i=1}^p \|y_i - x\|^2 |\ell_i(x)|$$

and

$$\|\nabla_x f(x) - \nabla_x m(x)\| \leq \kappa_{eg} \Lambda \Delta,$$

where $\Lambda = \max_{i=1, \dots, p} \max_{x \in \mathcal{B}(x, \Delta)} |\ell_i(x)|$.

[Ciarlet and Raviart, 1972]

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A simple DFO trust-region algorithm

- Compute an initial poised interpolation set Y_0
- Test for convergence
- **Build a quadratic model $m_k(x_k + s)$ of the objective function around an iterate x_k**

$$m_k(x_k + s) = f(x_k) + g(x_k)^T s + \frac{1}{2} s^T H s$$

based on well-poised sample sets.

- **Calculate a new trial point x_k^+ by solving**

$$\min_{s \in B(x_k; \Delta_k)} m_k(x_k + s).$$

in the trust region $B(x_k; \Delta_k)$.

A simple DFO trust-region algorithm

- Evaluate $f(x_k^+)$ and **compute the ratio**

$$\rho_k = \frac{f(x_k) - f(x_k + s_k)}{m(x_k) - m(x_k + s_k)} = \frac{\text{achieved reduction}}{\text{predicted reduction}}$$

- **Define the next iterate**
 - **case 1)** Successful iteration: set $x_{k+1} = x_k^+$, increase Δ_k and include point in the set Y_{k+1}
 - **case 2)** Unsuccessful iteration: set $x_{k+1} = x_k$, decrease Δ_k and include point in the set if its closer to x_k than the furthest in Y_k
- Compute the **new interpolation model m_{k+1}** around x_{k+1} using interpolation set Y_{k+1} if $Y_{k+1} \neq Y_k$, increment k

Geometry improving steps

- Fasano, Nocedal, and Morales [2009] observed that an algorithm which simply **ignores the geometry considerations** may in fact perform quite well in practice.
- But it may **lose** the property of **provable global convergence** to first-order critical points [Scheinberg and Toint, 2009].
- **Failure** of current iteration might be due to a too large trust region or a **bad quality** of the interpolation model (set not well-poised).
- Shows that we cannot afford to do without a geometry phase (need to maintain quality of the geometry of the interpolation set).
- Improvement is usually carried out at special "**geometry improving**" steps by computing additional function values at well-chosen points.

A DFO trust-region alg. with geometry restoration

- Compute an initial poised interpolation set Y_0
- Test for convergence
- Build a quadratic model $m_k(x_k + s)$ of the objective function around an iterate x_k

$$m_k(x_k + s) = f(x_k) + g(x_k)^T s + \frac{1}{2} s^T H s$$

based on well-poised sample sets.

- Calculate a new trial point x_k^+ by solving

$$\min_{s \in B(x_k; \Delta_k)} m_k(x_k + s).$$

in the trust region $B(x_k; \Delta_k)$.

A DFO trust-region alg. with geometry restoration

- Evaluate $f(x_k^+)$ and compute the ratio

$$\rho_k = \frac{f(x_k) - f(x_k + s_k)}{m(x_k) - m(x_k + s_k)}.$$

- Define the next iterate
 - **case 1)** Successful iteration: set $x_{k+1} = x_k^+$, increase Δ_k and include point in the set Y_{k+1}
 - **case 2)** Unsuccessful iteration: set $x_{k+1} = x_k$, decrease Δ_k and include point in the set if its closer to x_k than the furthest in Y_k
- Improve interpolation set by a geometry improving step
- Compute the new interpolation model m_{k+1} around x_{k+1} using interpolation set Y_{k+1} if $Y_{k+1} \neq Y_k$, increment k

Geometry improving steps

- As those geometry restoration steps are **expensive**, one may ask if they are really necessary.
- Idea is now to **reduce** the frequency and cost of the necessary tests as much as possible, while maintaining a mechanism for taking geometry into account.
- Design and convergence properties of new algorithm depend on a **self-correction mechanism** combining trust-region mechanism with polynomial interpolation setting.

The new DFO trust-region algorithm

- Compute an initial poised interpolation set Y_0
- **Test for convergence and improve geometry if necessary**
- **Build a quadratic model $m_k(x_k + s)$** of the objective function around an iterate x_k

$$m_k(x_k + s) = f(x_k) + g(x_k)^T s + \frac{1}{2} s^T H s$$

based on the current interpolation set.

- **Calculate a new trial point x_k^+** by solving

$$\min_{s \in B(x_k; \Delta_k)} m_k(x_k + s).$$

in the trust region $B(x_k; \Delta_k)$.

The new DFO trust-region algorithm

- Evaluate $f(x_k^+)$ and compute the ratio ρ_k
- Define the next iterate
 - **case 1)** Successful iteration: include point in the set Y_{k+1} , adjust Δ and define $x_{k+1} = x_k^+$
 - **case 2)** Try to replace a far interpolation point: if set F_k is non-empty, include point in the set Y_{k+1} , set $\Delta_{k+1} = \Delta_k$
 - **case 3)** Try to replace a close interpolation point: if set $F_k = \emptyset$ and set C_k is non-empty, include point in the set Y_{k+1} , set $\Delta_{k+1} = \Delta_k$
 - **case 4)** Reduce trust-region radius and set $Y_{k+1} = Y_k$.

Self-correcting geometry

Set of **far** points:

$$F_k = \{y_{k,j} \in Y_k \text{ such that } \|y_{k,j} - x_{best}\| > \beta\Delta \text{ and } \ell_{k,j}(x_k^+) \neq 0\}$$

Set of **close** points:

$$C_k = \{y_{k,j} \in Y_k \text{ such that } \|y_{k,j} - x_{best}\| \leq \beta\Delta \text{ and } \ell_{k,j}(x_k^+) > \Lambda\}$$

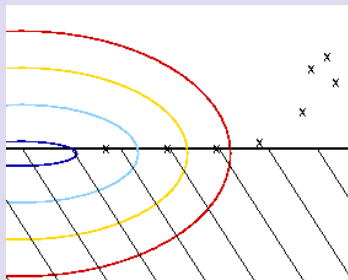
Self-correcting property:

If iteration k is unsuccessful, $F_k = \emptyset$ and $\Delta_k \leq \kappa_\Lambda \|\nabla m_k\|$, then $C_k \neq \emptyset$, and so, every unsuccessful iteration must result in an improvement of the interpolation set geometry. [Scheinberg and Toint, 2009]

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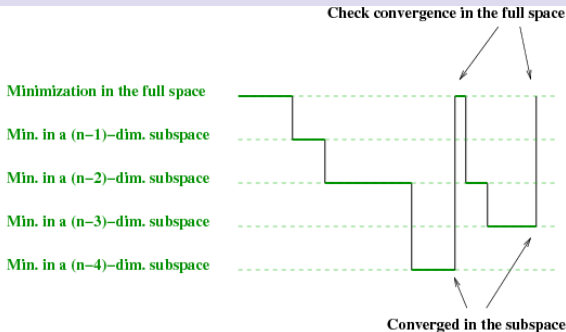
Extension to bounds



- **Situation:** algorithm converges towards a minimum
- **Problem:** iterates get aligned along the bound
- Results in a **degenerate set of points** due to the bounds!
- Λ -poisedness **no suitable measure** anymore, because maximum of Lagrange polynomials lies outside of the bounds
- Thus, self correcting property **not working**

Solution: a subspace method

Continue minimization in a smaller dimensional subspace



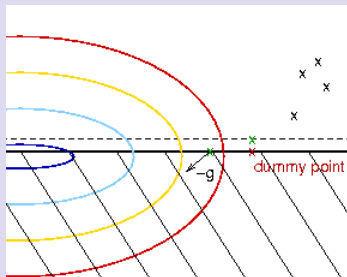
- If encounter an active bound, **reduce dimensionality**
- If converged in the subspace, going back to **check convergence** in the full space

Further features of the algorithm (I)

- Degree of **initial interpolation model** user-defined: linear, diagonal, quadratic
- **Adjust** the initial trust region and **shift** the starting point to build the initial model **inside the bounds**
- Using variable size models: unless model is quadratic, new **iterates augment the size** of the interpolation set
- Initial degree of subspace-models is linear and is then augmented with the new iterates computed in the subspace
- **Recursive technique**: call the algorithm itself to solve the problem in the subspace(s)

Further features of the algorithm (II)

Attempt to save function evaluations by creating dummy points



- New active bound: need to build a model in the subspace
- Consider points lying **close** to the active bound(s), create **dummy points**
- **Compute the model values** at the dummy points
- Take real points and dummy points lying in the subspace to **build the model**
- Dummy points are then **replaced** by the new iterates

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Methodology

CUTer testing environment

- 50 bound-constrained test cases from CUTer test environment
- Nbr. of variables varies from 1 to 25 dimensions

Competitor: BOBYQA

- State of the art software developed by M.J.D. Powell [2006]
- Currently one of the best codes for bound-constrained minimization without derivatives

Stopping criterion

- Stopping criteria are different
- Using optimal objective function value computed by TRON (using first and second derivatives) as a reference
- We terminate when 6 correct significant figures in f were attained

Numerical results

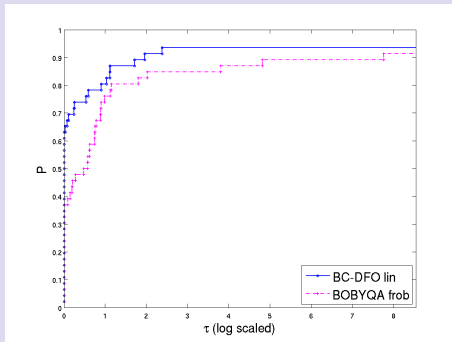


Figure: Performance profile in terms of nbr. of function eval.

A success in solving a 25-dim. problem

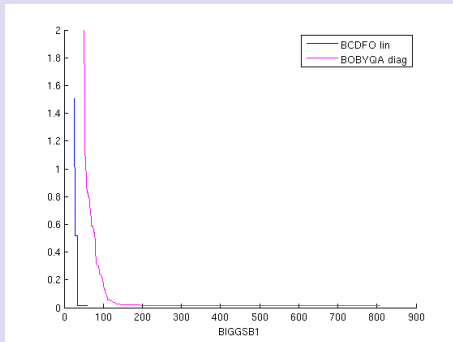


Figure: Convergence history of problem BIGGSB1

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Conclusions and Perspectives

Summary

- Presented a new model-based trust-region DFO algorithm with a self-correcting geometry property
- Extended the algorithm to handle bounds
- Implemented a robust version of the algorithm: BC-DFO
- Compared BC-DFO to BOBYQA with quite satisfying results

Perspectives

- Consider further enhancements on model Hessian update to improve performance
- Test the algorithm on real-life application (aerodynamic functions provided by Airbus)
- Implement the use of an inexact gradient

Thank you for your attention!