

Towards an optimal-order approximate sparse factorization exploiting data-sparseness in separators

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Sparse Days, June 17-18, CERFACS

Motivation: sparse factorization

- ✓ Numerically robust, efficient for many RHS.
- ✗ Expensive. For model problems on regular grids (mesh size k):
 - 2D ($N = k^2$): flops = $\Theta(N \log N)$, memory = $\Theta(N^{3/2})$.
 - 3D ($N = k^3$): flops = $\Theta(N^{4/3})$, memory = $\Theta(N^2)$.

Example: finite difference, $1000 \times 1000 \times 1000$ grid, 27pt stencil:
flops=50 exaflops, mem>200 TB (the whole Cray XE6 system at NERSC!).

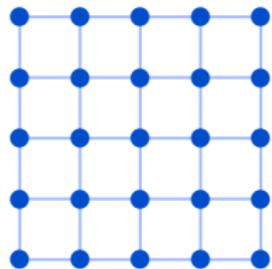
Idea

Resort to approximation.

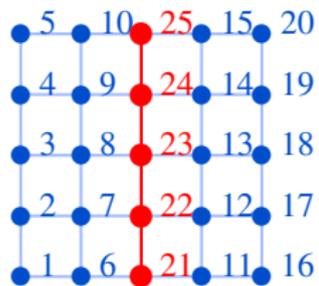
Outline

- Review of multifrontal LU
- Structured partial factorization: HSS-embedded multifrontal
- Parallelization, performance
- Remarks

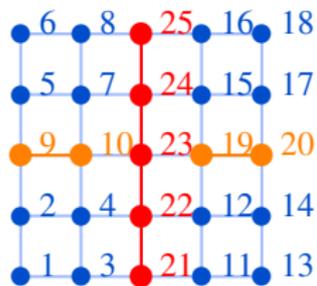
Nested dissection ordering



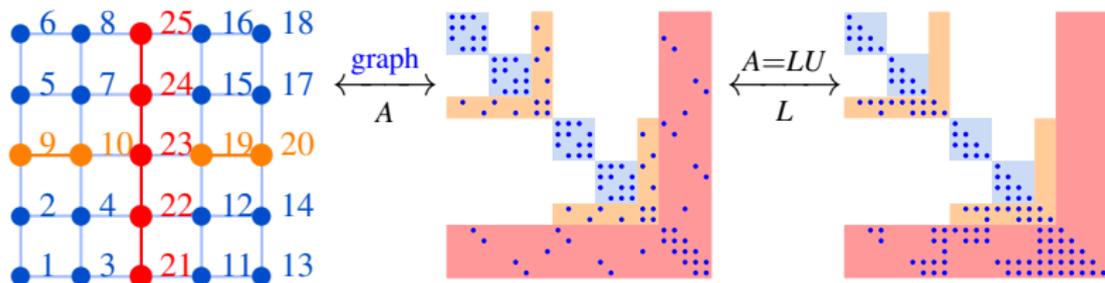
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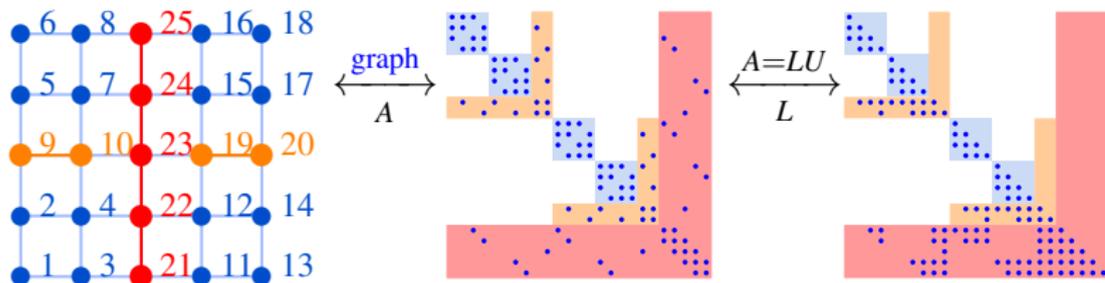
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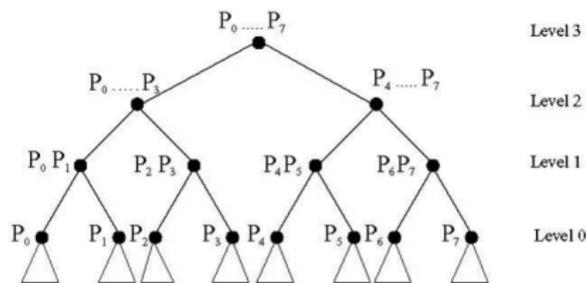
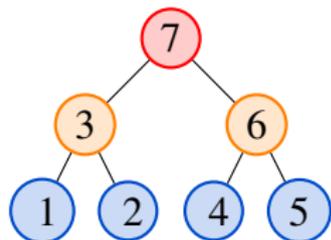
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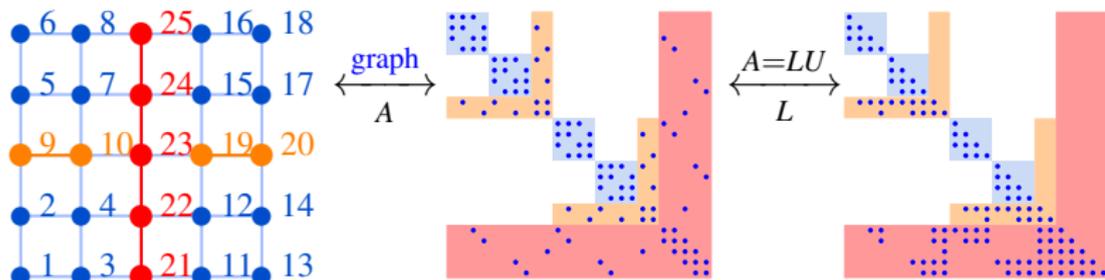
Nested dissection ordering



- Better parallel properties

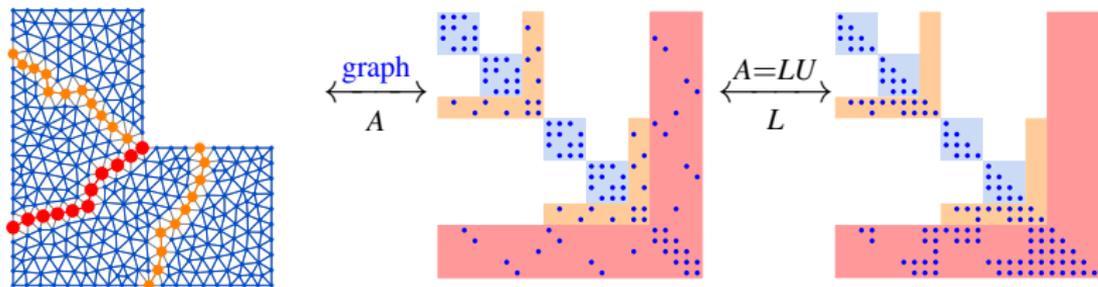


Nested dissection ordering



- Better parallel properties
- Fewer entries in factors (here, 129 vs. 116)
- Implementation: using [Scotch](#) partitioner, since separator info needed

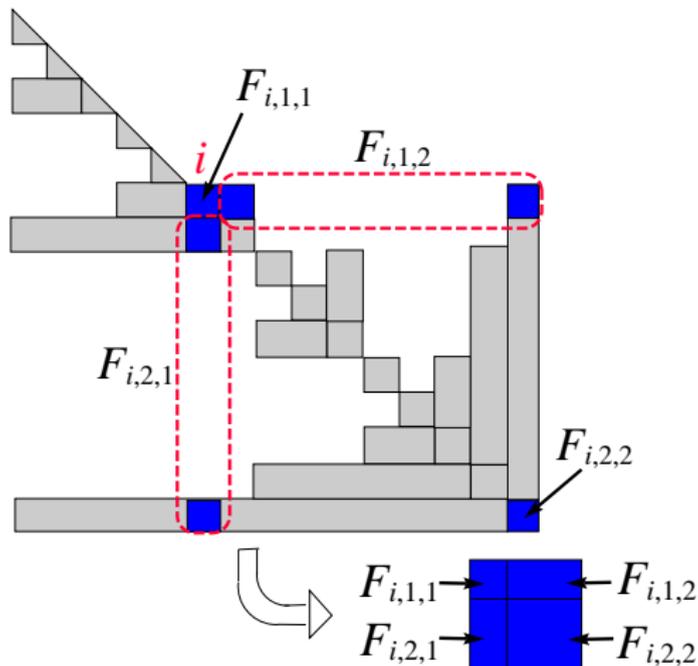
Nested dissection ordering



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Frontal and Update matrices

- Each separator corresp. to a **dense** submatrix (frontal & update)
 - ▶ **Dominant cost:** $F_{1,1} = LU, F_{2,1}, F_{1,2}$, Update matrix



Observation : presence of data-sparseness

Several approaches are available in the literature that exploit low-rank in dense / structured matrices:

- Hierarchical matrices: \mathcal{H} , \mathcal{H}^2 , (Bebendorf, Börm, Grasedyck, Hackbusch, ...)
- Hierarchically/Sequential Semiseparable (HSS / SSS) representations (Chandrasekaran, Dewilde, Gu, Li, Olshevsky, Vandebril, Xia, ...)
- BLR (Amestoy et al.)

Some representations are simpler and apply to broader classes of problems but provide less gain in memory/operations;

Some others are more complex but allow for further gains in complexity.

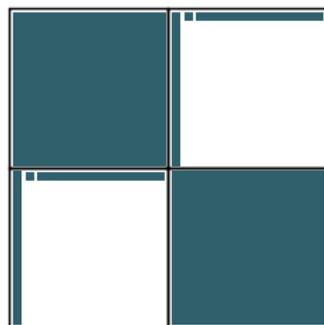
We focus on Hierarchically Semi-Separable matrices (HSS) ...

Hierarchically Semi-Separable matrices

An HSS matrix A is a dense matrix whose **off-diagonal blocks are low-rank**.

High-level structure: 2×2 blocks

$$A = \left[\begin{array}{c|c} D_1 & U_1 B_1 V_1^T \\ \hline U_2 B_2 V_1^T & D_2 \end{array} \right]$$



Fundamental property required for efficiency: **nested bases**

$$U_3 = \begin{bmatrix} U_1 & 0 \\ 0 & U_2 \end{bmatrix} U_3^{small}, U_3^{small} : 2k \times k$$

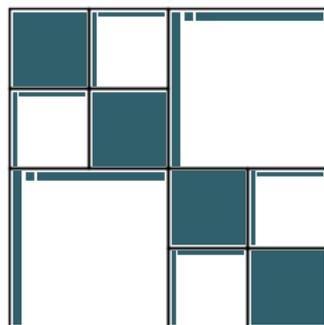
Same for U_3, U_6, V_6 and recursively at subsequent levels.

Hierarchically Semi-Separable matrices

An HSS matrix A is a dense matrix whose **off-diagonal blocks are low-rank**.

Recursion

$$A = \left[\begin{array}{cc|c} D_1 & U_1 B_1 V_2^T & U_3 B_3 V_6^T \\ \hline U_2 B_2 V_1^T & D_2 & \\ \hline U_6 B_6 V_3^T & D_4 & U_4 B_4 V_5^T \\ \hline & U_5 B_5 V_4^T & D_5 \end{array} \right]$$



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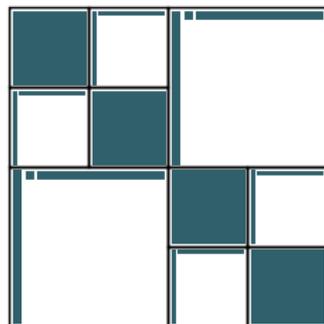
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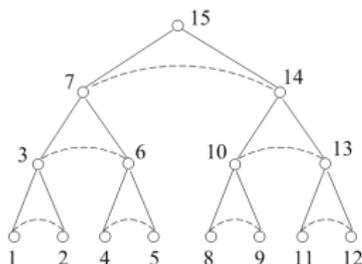
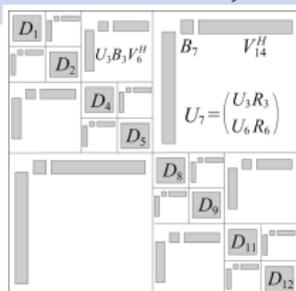


Fundamental property required for efficiency: **nested bases**

$$U_3 = \begin{bmatrix} U_1 & 0 \\ 0 & U_2 \end{bmatrix} U_3^{small}, U_3^{small} : 2k \times k$$

Same for U_3 , U_6 , V_6 and recursively at subsequent levels.

Hierarchical bases, HSS tree



For efficiency, require:

$$U_3 = \begin{bmatrix} U_1 & 0 \\ 0 & U_2 \end{bmatrix} U_3^{small}, \quad U_3^{small} : 2k \times k, \quad U_6 = \begin{bmatrix} U_4 & 0 \\ 0 & U_5 \end{bmatrix} U_6^{small}, \quad U_6^{small} : 2k \times k$$

$$U_7 = \begin{bmatrix} U_3 & 0 \\ 0 & U_6 \end{bmatrix} U_7^{small}, \quad U_7^{small} : 2k \times k$$

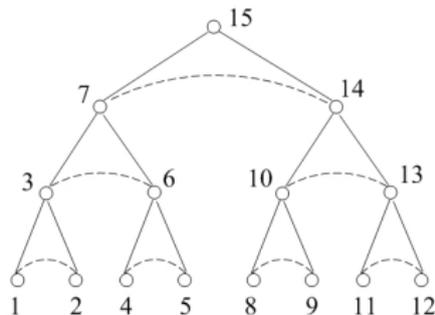
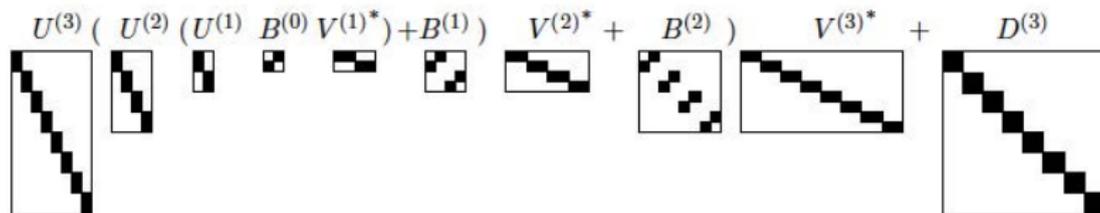
Each basis is a product of descendants' bases:

$$U_7 = \begin{bmatrix} U_1 & 0 & 0 & 0 \\ 0 & U_2 & 0 & 0 \\ 0 & 0 & U_4 & 0 \\ 0 & 0 & 0 & U_5 \end{bmatrix} \begin{bmatrix} U_3^{small} & 0 \\ 0 & U_6^{small} \end{bmatrix} U_7^{small},$$

Not to multiply out!

HSS explicit representation (construction)

[Martinsson]



- keep it as an unevaluated product & sum
- operations going up / down the HSS tree

HSS operations

With r the **maximum rank** of a block:

- HSS construction (compression of low-rank blocks): $\Theta(rN^2)$.
- ULV factorization: $\Theta(r^2N)$ [Chandrasekaran et al.].
- HSS solution $\Theta(rN)$.

Provable r for some discretized PDEs: [Chandrasekaran et al.'10, Enquist-Ying'11]

2D Poisson	$\Theta(1)$
2D Helmholtz	$\Theta(\log k)$
3D Poisson	$\Theta(k)$
3D Helmholtz	$\Theta(k)$

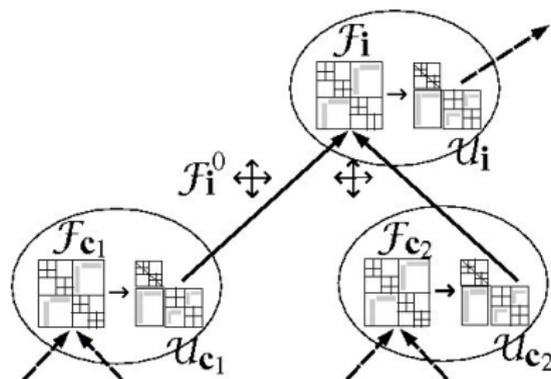
Constructing the HSS structure is the dominant operation.

Embedding HSS in multifrontal

Approximate **Frontal** & **Update** matrices by HSS

Operations:

- HSS construction of frontal \mathcal{F}_i , ULV factorize $\mathcal{F}_i(1, 1)$
- HSS approximation of $\mathcal{U}_i, \mathcal{U}_j$
- extend-add of HSS matrices \mathcal{U}_i and \mathcal{U}_j to parent



Difficulty: extend-add of two non-matching HSS structures

Embedding HSS in multifrontal

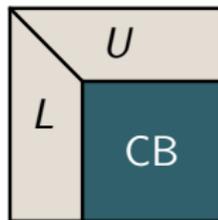
Several ways to deal with frontal matrices:

More complicated \uparrow
More memory \downarrow

Fully structured: HSS on the whole frontal matrix. No dense matrix.

Partial+: HSS on the whole frontal matrix. Dense frontal matrix.

Partially structured: HSS on the L, U part only. Dense frontal matrix, dense CB in stack after partial factorization.



Once a frontal matrix is in partial or complete HSS form, an **ULV factorization** is applied instead of usual LU factorization.

Complexity (for fully-structured) [Xia '11, Chandrasekaran et al. '11]:

Problem	Classical MF		MF with HSS	
	Mem	Flops	Mem	Flops
2D Poisson	$\Theta(N \log N)$	$\Theta(N^{3/2})$	$\Theta(N \log \log N)$	$\Theta(N \log N)$
2D Helmholtz	$\Theta(N \log N)$	$\Theta(N^{3/2})$	$\Theta(N \log \log N)$	$\Theta(N \log N)$
3D Poisson	$\Theta(N^{4/3})$	$\Theta(N^2)$	$\Theta(N \log N)$	$\Theta(N^{4/3})$
3D Helmholtz	$\Theta(N^{4/3})$	$\Theta(N^2)$	$\Theta(N \log N)$	$\Theta(N^{4/3} \log N)$

Experiments

- Hopper, Cray XE6 at NERSC.
- Hsolver (for Helmholtz), geometric HSS-embedded multifrontal solver
- Helmholtz equations with PML boundary

$$\left(-\Delta - \frac{\omega}{v(x)^2}\right) u(x, \omega) = s(x, \omega)$$

Δ : Laplacian

ω : angular frequency

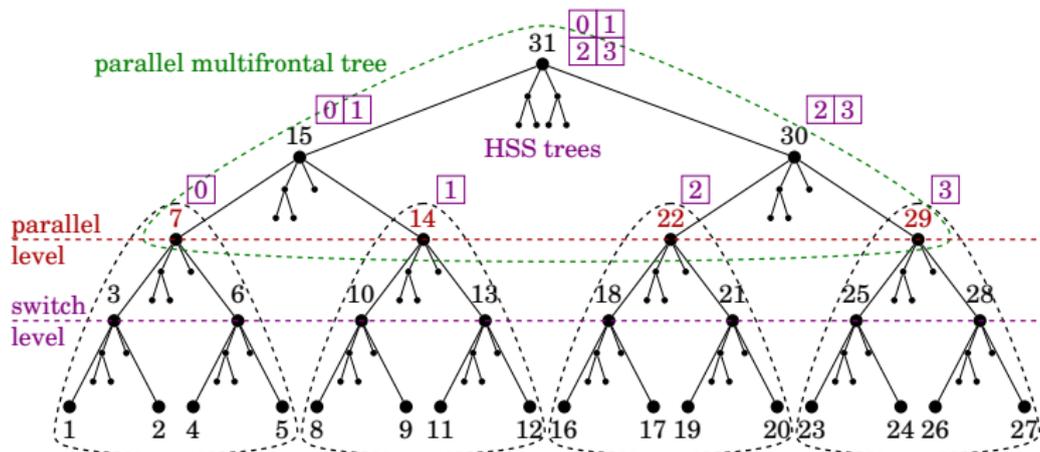
$v(x)$: seismic velocity field

$u(x, \omega)$: time-harmonic wavefield equation

- Finite Difference discretized system: complex, pattern-symmetric, non-Hermitian. **Indefinite, ill-conditioned**
- single precision

MF + HSS: two types of tree-based parallelism

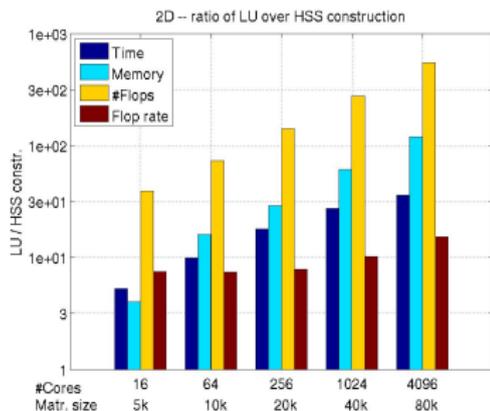
- **Outer tree:** separator tree for multifrontal factorization
- **Inner tree:** HSS tree at each internal separator node



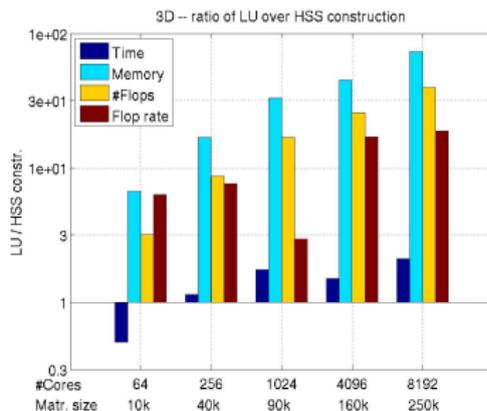
HSS performance

- HSS construction on the last Schur complement corresp. to the top separator.

Performance ratio of LU over HSS:



(a) 2D, max_rank=7



(b) 3D, max_rank=848

Sparse results - 2D problems

2D Helmholtz problems on square grids (mesh size k , $N = k^2$), 10 Hz.

k		10,000	20,000	40,000	80,000
P		64	256	1,024	4,096
MF	Factorization (s)	258.6	544.8	1175.8	2288.5
	Gflops/s	507.3	2109.3	8185.6	31706.9
	Solution+refinement (s)	10.4	10.8	11.5	11.6
	Factors size (GB)	120.1	526.7	2291.2	9903.7
	Max. peak (GB)	2.3	2.5	2.7	2.9
	Communication volume (GB)	136.2	1202.5	9908.1	
HSS	HSS+ULV (s)	97.9	172.5	325.3	659.3
	Gflops/s	196.9	715.6	2820.7	9820.6
	Solution+refinement (s)	20.2	55.4	61.4	115.8
	Steps	3	3	9	9
	Factors size (GB)	66.2	267.7	1333.2	4572.3
	Max. peak (GB)	1.7	1.7	1.7	1.7
	Communication volume (GB)	74.2	573.8	4393.4	
	HSS rank	258	503	1013	2015
$\ x - x_{MF}\ / \ x_{MF}\ $		1.5×10^{-5}	2.2×10^{-5}	3.1×10^{-5}	3.5×10^{-6}
$\max_i \frac{ Ax-b _i}{(A x + b)_i}$		7.1×10^{-6}	1.0×10^{-5}	2.0×10^{-6}	3.5×10^{-6}

Results - 3D problems

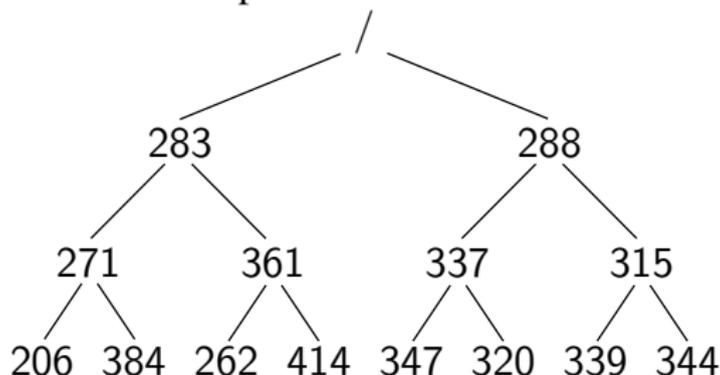
3D Helmholtz problems on cubic grids (mesh size k , $N = k^3$), 5 Hz.

k		100	200	300	400
P		64	256	1,024	4,096
MF	Factorization (s)	88.4	1528.0	1175.8	6371.6
	Gflops/s	600.6	2275.7	9505.6	35477.3
	Solution+refinement (s)	0.6	2.2	3.5	4.8
	Factors size (GB)	16.6	280.0	1450.1	4636.1
	Max. peak (GB)	0.5	1.9	2.5	2.0
	Communication volume (GB)	83.1	2724.7	26867.8	165299.3
HSS	HSS+ULV (s)	120.4	1061.3	2233.8	3676.5
	Gflops/s	207.8	720.4	2576.6	6494.8
	Solution+refinement (s)	2.3	8.2	31.5	182.8
	Steps	4	5	10	6
	Factors size (GB)	10.7	112.9	434.3	845.3
	Max. peak (GB)	0.5	1.7	2.1	0.4
	Communication volume (GB)	93.6	2241.2	18621.1	143300.0
	HSS rank	481	925	1391	1860
$\ x - x_{MF}\ / \ x_{MF}\ $		6.2×10^{-6}	9.4×10^{-7}	1.1×10^{-6}	1.7×10^{-6}
$\max_i \frac{ Ax-b _i}{(A x + b)_i}$		1.5×10^{-7}	5.7×10^{-7}	9.7×10^{-7}	3.7×10^{-6}

Performance analysis – rank imbalance

- Rank revealing QR is the dominant operation (2/3 of the total time).
- Observed some load imbalance (factors up to 2.5) due to some imbalance in the ranks found in frontal matrices (run time of RRQR is proportional to the rank).

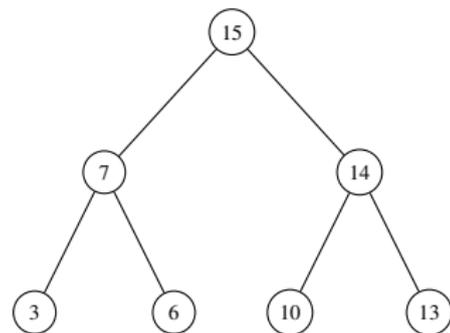
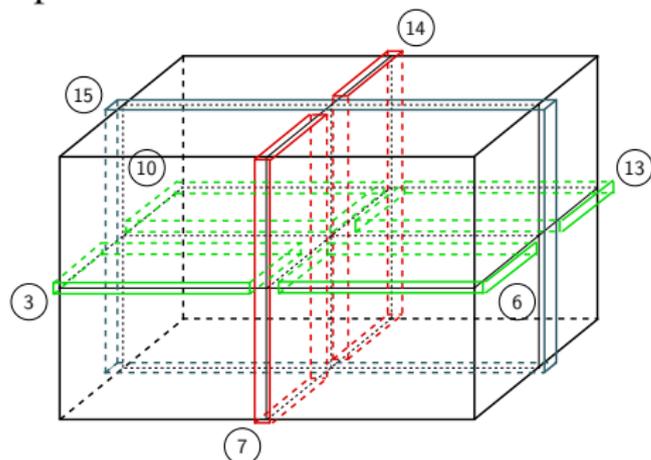
Example: ranks in the row compression of the root of a 100^3 problem



HSS tree; 8 leaves/8 processes.

Ordering of separators/fully-summed variables

In a regular 3D mesh with nested dissection, each frontal matrix corresponds to a plane:

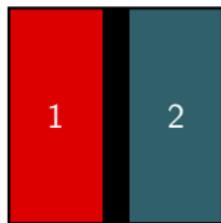


Top of the elimination tree.

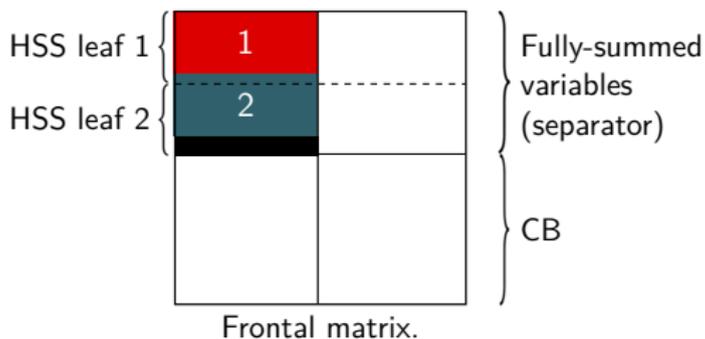
Within a frontal matrix/separator, the ordering of the variables is crucial to get low rank blocks (cf. [admissibility condition](#) for PDEs [Börm, Grasedyck, Hackbush]).

Ordering of separators/fully-summed variables – 2

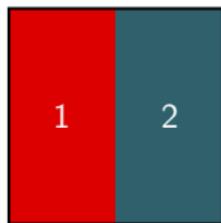
Vertex-based approach:



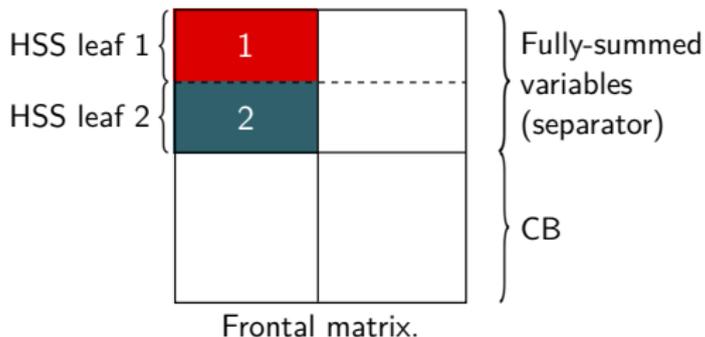
2D separator from a 3D domain.



Edge-based approach:



2D separator from a 3D domain.



- An edge-based Nested Dissection / Morton ordering better preserves geometry and should provide better balance in ranks.
- On some medium-sized problems we observed $\sim 10\%$ gain in run time. [Work in progress.](#)

Summary, future

- Parallel geometric HSS-embedded MF solver for Helmholtz equations is faster than a regular MF solver. **Gains increase with problem size.**
- Using iterative refinement, it delivers accurate solutions.

- Explore ways to reduce further the memory footprint, in particular the stack of **contribution blocks**. Structured Schur complement?
Randomized sampling?
- Move towards a parallel **algebraic** solver
- Analyze communication bound
- Black-box preconditioner? (compare to ILU, etc.)
- Compare to sparse solvers using other low-rank forms [Saad et al., Weisbecker et al., Ying et al.]
- Resilience at extreme scale