## Towards a Scalable Parallel Sparse Linear System Solver

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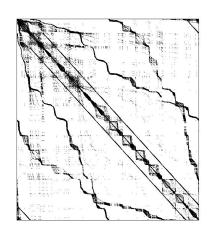
#### **Outline**

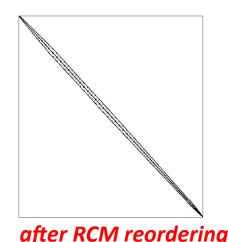
- Motivating applications a sample
- The parallel scalability of the SPIKE family of banded linear system solvers
- Extensions for solving general sparse linear systems:
  - The Pardiso-SPIKE hybrid solver (PSPIKE)
- Role of reordering for faster MATVEC and extraction of preconditioners:
  - TraceMIN a parallel eigensolver for computing the Fiedler vector (spectral reordering)

## Target Computational Loop

| Integration           |
|-----------------------|
| Newton iteration      |
| Linear system solvers |
| $oldsymbol{\eta}_k$   |
| $\leftarrow_{k}$      |
| -                     |

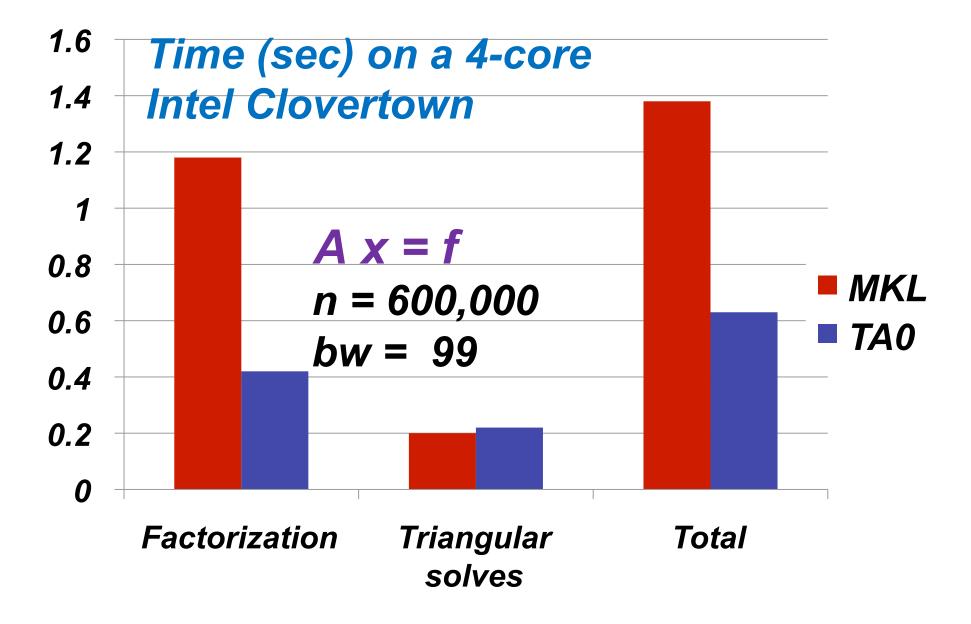
### Role of Reordering Schemes

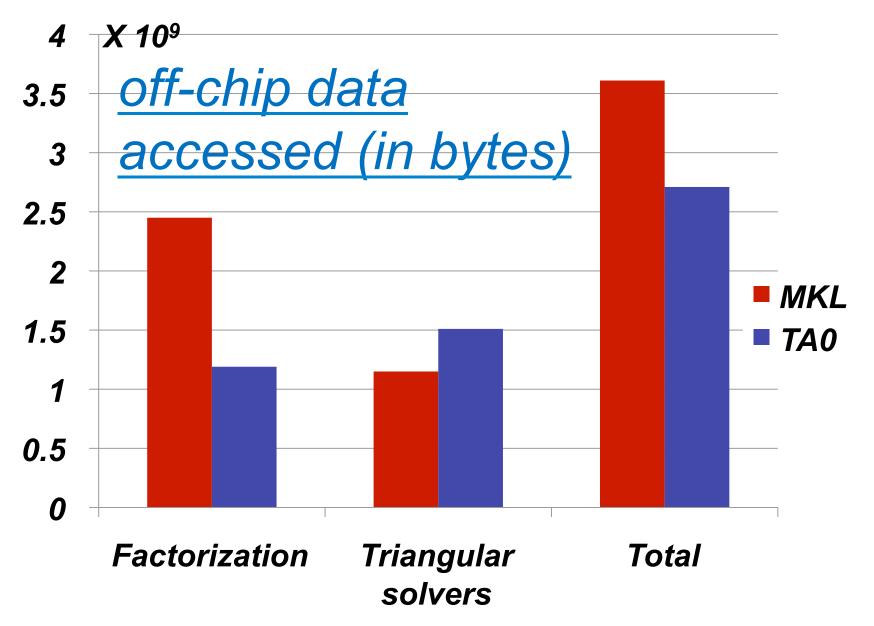




<u>Banded</u> <u>Preconditioners</u>

- "Banded", or low-rank perturbations of banded, systems are often obtained after RCM or spectral reordering,
- This yields more scalable parallel matrix-vector multiplications, and helps in extracting more effective "banded" preconditioners.





"Analyzing memory access intensity in parallel programs for multicore architectures"

L. Liu, Z. Li, and A. S.

#### Spike-based sparse solver: PSPIKE

Step 1: reordering (equivalent to HSL -- MC64 + MC73)

Step 2: extract a "banded" preconditioner M



(or any Krylov subspace method)

Solve Mz = r (via Pardiso-SPIKE)

Manguoglu M, Sameh A, and Schenk O, PSPIKE: A Parallel Hybrid Sparse Linear System Solver, Lecture Notes in Computer Science(proceedings of EuroPar09), Volume 5704, pp.797-808, 2009

# Weighted Spectral Reordering (WSO) Main Objectives:

- -encapsulating as many of the heaviest sparse matrix elements as possible within a central band to be used as a preconditioner,
  - reducing the rank of the matrix lying outside the central band.
- -realizing a faster "matvec"

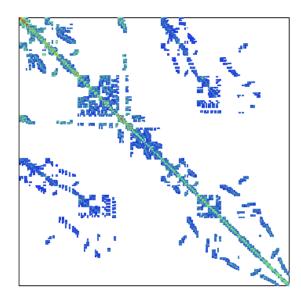
(banded matrix + few nonzero elements outside the band)

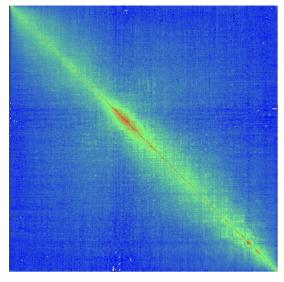
#### **UFL:** smt -- structural mechanics

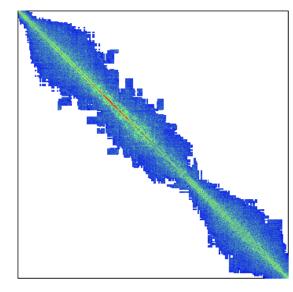
N: 25,710 NNZ: 3,749,582



after TraceMIN-Fiedler







Original matrix

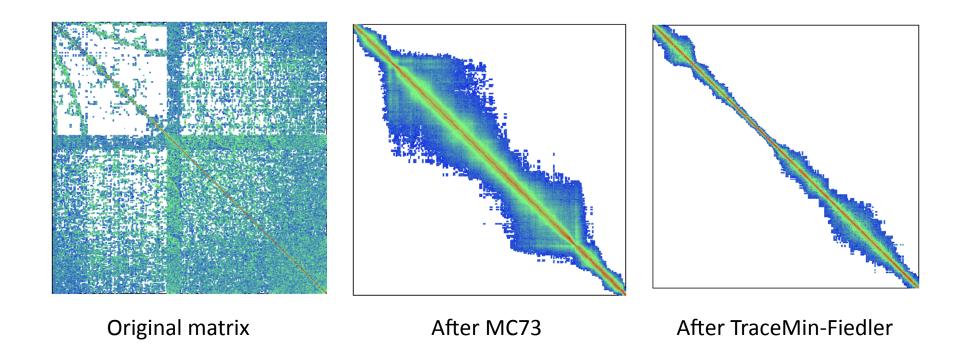
After MC73

After TraceMin-Fiedler

obtaining the Fiedler vector via the eigensolver: TraceMIN (Wisniewski and A.S. -- SINUM, '82)

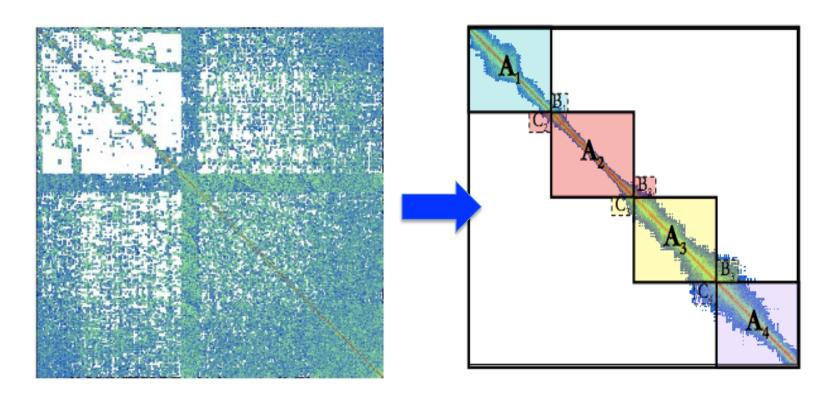
#### UFL: f2 -- structural mechanics

N: 71,505 NNZ: 5,294,285



TraceMIN-Fiedler: Murat Manguoglu et. al -- submitted

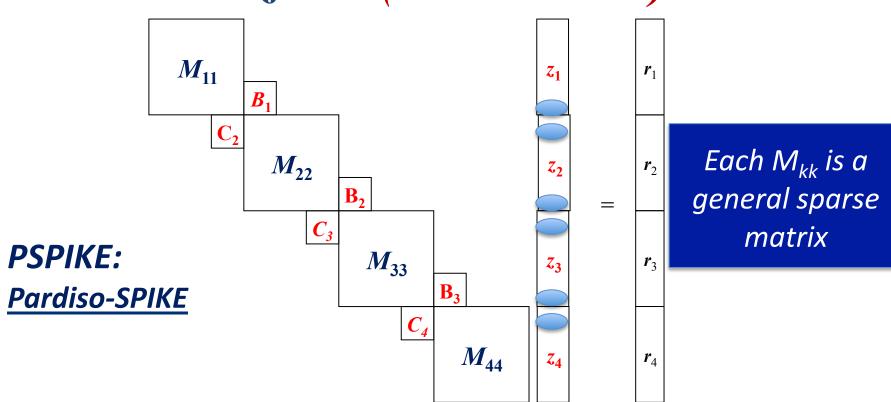
## UFL-f2



Before reordering

After reordering via TraceMIN-Fiedler

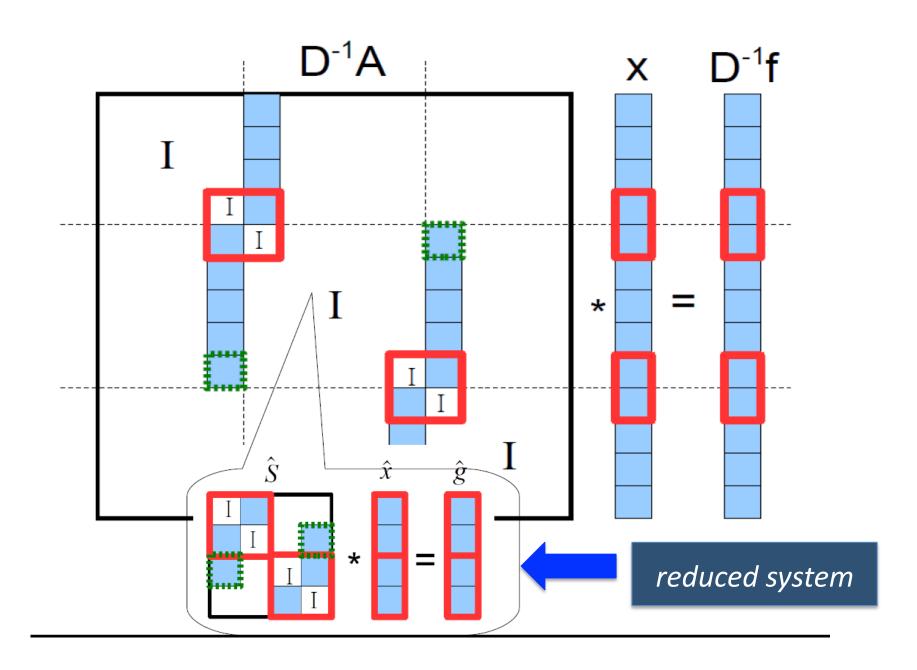
#### Mz = r (M is "banded")



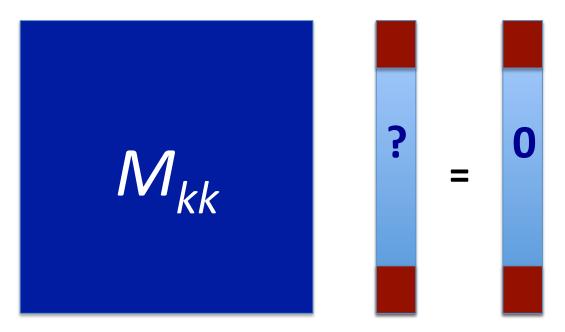
$$P = M + \delta(M) = D' * S'$$

- (i) Solve D'y = r

Solving systems involving (ii) Solve S'z = y The preconditioner Pz = r

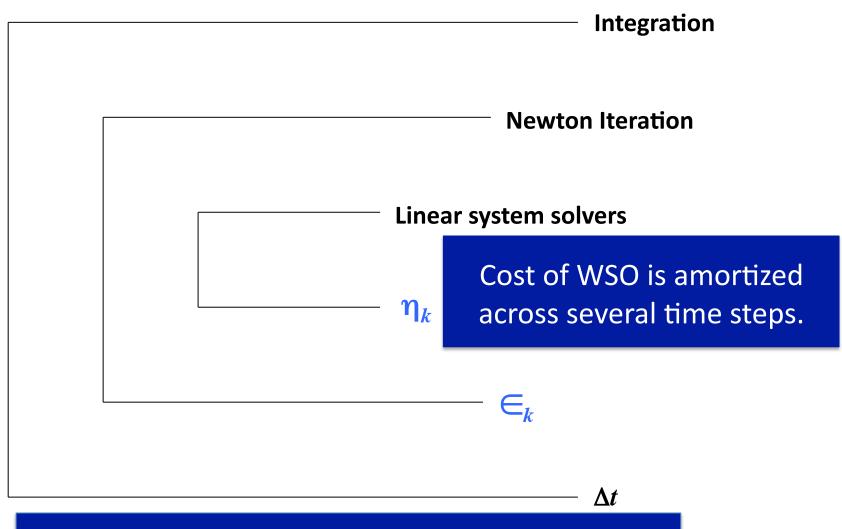


#### Generating tips of the spikes



Obtain the upper and lower tips of the solution block via the modified direct sparse system solver "Pardiso".

#### How expensive is spectral reordering?



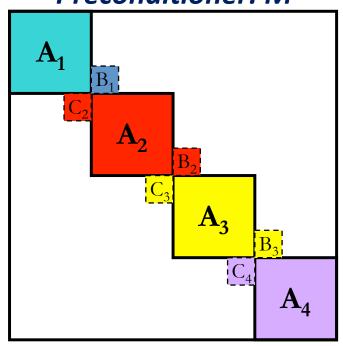
will return to this issue later

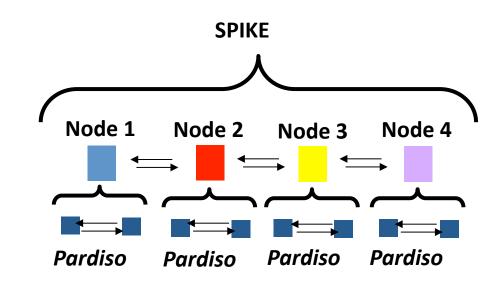
## $\sigma = T(MC73)$ on 1 core $\div$ T(TraceMIN-Fiedler) on 4 nodes

| Matrix          | n       | nnz    | σ   |
|-----------------|---------|--------|-----|
| Rajat31         | ~4.7 M  | ~ 20 M | 85  |
| Schenk/nlpkkt20 | ~ 3.5 M | ~ 95 M | 8   |
| Freescale1      | ~ 3.4 M | ~17 M  | 4   |
| KKT-power       | ~ 2.1 M | ~ 13 M | 641 |

## Multilevel Parallelism of PSPIKE for solving Mz = r

#### **Preconditioner: M**







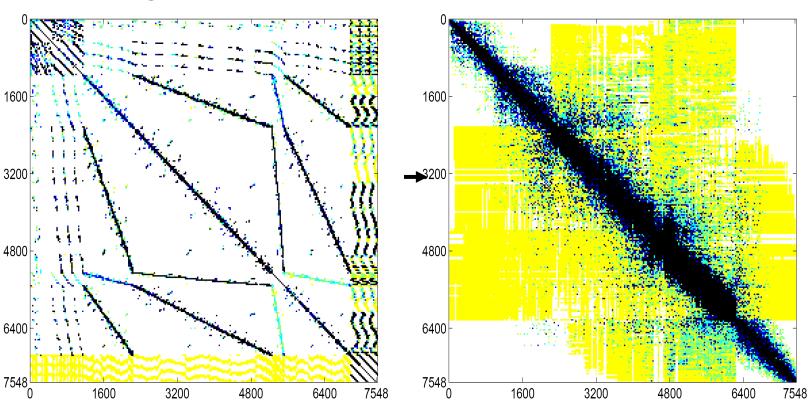
Preconditioner M is generalized-banded (MBP)
Other cases:

- narrow-banded (NBP)
- wide-banded (WBP)

#### **Circuit Simulation**

#### Original Matrix

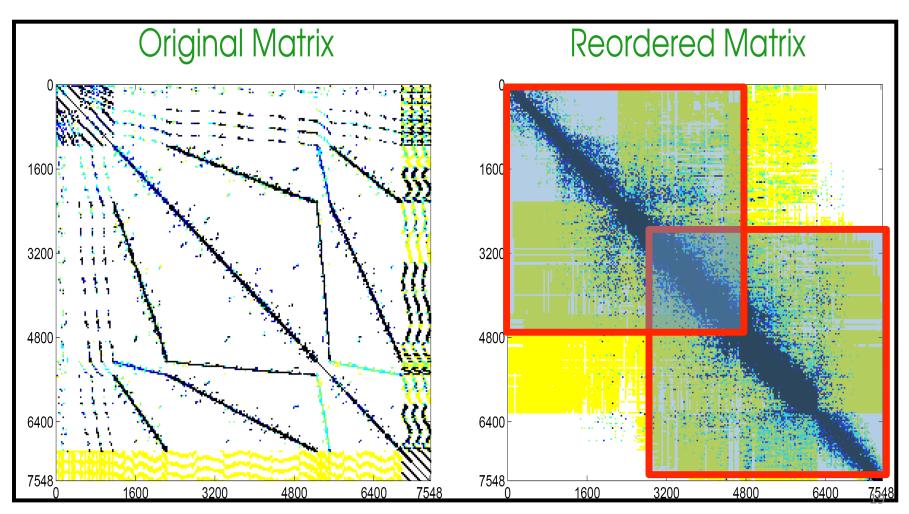
## after TraceMIN-Fiedler Reordered Matrix



Wide-Banded Preconditioner

#### WBP: wide-banded preconditioners

preconditioner consists of overlapped blocks Naumov, Manguoglu, A.S.: A Tearing-based Hybrid Parallel Sparse Linear System Solver: JACM, 2010.



#### Solve: Ax = f; A := nonsingular

$$\begin{bmatrix} A_{11} & A_{12} & A_{2} \\ A_{21} & A_{22} & A_{23} \\ A_{32} & A_{33} \end{bmatrix} \begin{bmatrix} x_{1} \\ x_{2} \\ x_{3} \end{bmatrix} = \begin{bmatrix} f_{1} \\ f_{2} \\ f_{3} \end{bmatrix}$$

$$\begin{bmatrix} A_{11}^{(1)} & A_{12}^{(1)} \\ A_{21}^{(1)} & A_{22}^{(1)} \end{bmatrix} \begin{bmatrix} x_{1} \\ x_{2} \end{bmatrix} = \begin{bmatrix} f_{1} \\ \alpha f_{2} + y \end{bmatrix}; \ 0 < \alpha < 1$$

$$(n \times n) \begin{bmatrix} A_{11}^{(2)} & A_{12}^{(2)} \\ A_{21}^{(2)} & A_{22}^{(2)} \end{bmatrix} \begin{bmatrix} \hat{x}_{2} \\ x_{3} \end{bmatrix} = \begin{bmatrix} (1 - \alpha) f_{2} - y \\ f_{3} \end{bmatrix}$$

$$(n \times n) \begin{bmatrix} A_{12}^{(2)} & A_{22}^{(2)} \\ A_{21}^{(2)} & A_{22}^{(2)} \end{bmatrix} \begin{bmatrix} \hat{x}_{2} \\ x_{3} \end{bmatrix} = \begin{bmatrix} (1 - \alpha) f_{2} - y \\ f_{3} \end{bmatrix}$$

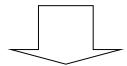
$$A_{22} = A_{22}^{(1)} + A_{11}^{(2)} \in \Re^{m \times m}$$

$$m < < n$$

$$A_{1} = \begin{bmatrix} A_{11}^{(1)} & A_{12}^{(1)} \\ A_{21}^{(1)} & A_{22}^{(1)} \end{bmatrix} \begin{bmatrix} x_{1} \\ x_{2} \end{bmatrix} = \begin{bmatrix} f_{1} \\ \alpha f_{2} + y \end{bmatrix}$$

$$A_{2} = \begin{bmatrix} A_{11}^{(2)} & A_{12}^{(2)} \\ A_{21}^{(2)} & A_{22}^{(2)} \end{bmatrix} \begin{bmatrix} \hat{x}_{2} \\ x_{3} \end{bmatrix} = \begin{bmatrix} (1 - \alpha) f_{2} - y \\ f_{3} \end{bmatrix}$$

Determine that  $\dot{y}$  which assures that  $x_2 = \hat{x}_2$ .



 $\dot{y}$  is the solution of the *balance system* (order = size of overlap)

$$My = g$$

M and g are not available explicitly

### The Balance System

- Solve the system My=g using a Krylov subspace scheme (CG or Bicgstab).
- Two time-consuming kernels:

$$-q = M*p$$

$$-r(p)=g-M*p$$

## **Observations**

• 
$$r(p) = g - M*p$$
  
=  $\hat{x}_2(p) - x_2(p)$ 

• 
$$r(0) = g$$
  
=  $\hat{x}_2(0) - x_2(0)$ 

• 
$$q = M*p$$
  
=  $r(0) - r(p)$ 

### **PSPIKE**

#### Stage 1 –

- Extraction of an effective "banded" preconditioner
   M with or without reordering.
- $norm(M,'fro') = (1 \varepsilon) norm(A,'fro');$
- "bandwidth"  $\leq \beta(n)$ .

#### Stage 2 –

- Use an outer Krylov subspace method (BiCGstab) to solve A x = f with preconditioner M.
- A modified Pardiso (by O. Schenk) is a major kernel for solving Mz = r.
- Classes of banded preconditioners:

narrow-banded, generalized-banded, wide-banded

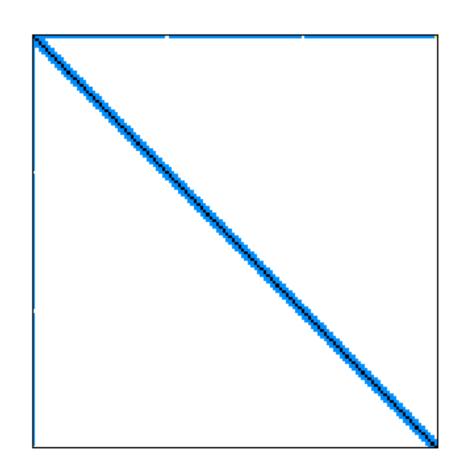
## Computing Platform

- Intel Cluster
  - -Westmere X5670, 2.93 GHz
  - —Infiniband interconnection
  - -12-core nodes

## a. PSPIKE-NBP

#### UFL - Rajat31 (circuit simulation)

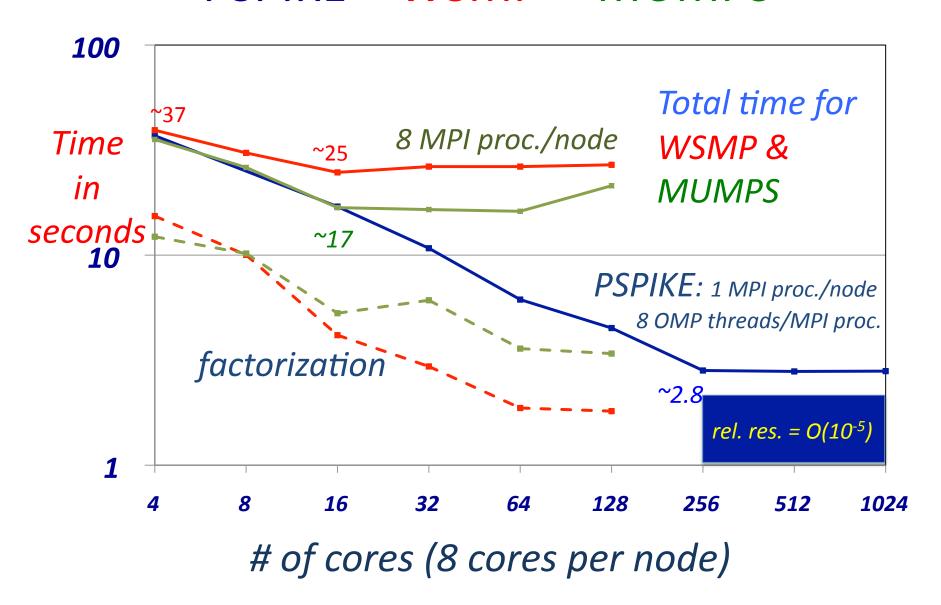
N ~ 4.7 M
nnz ~ 20 M
nonsymmetric



## Scalability of PSPIKE vs. direct solvers (for Rajat31)

- Direct solvers:
  - WSMP
    - matrix input centralized on Node 0
  - MUMPS
    - matrix input distributed by block rows
    - solution vector distributed
    - ParMETIS
- *PSPIKE*: (bandwidth of preconditioner = 5)
  - TraceMIN-Fiedler
  - matrix input distributed by block rows

#### PSPIKE - WSMP -- MUMPS



#### MEMS simulation benchmark 1

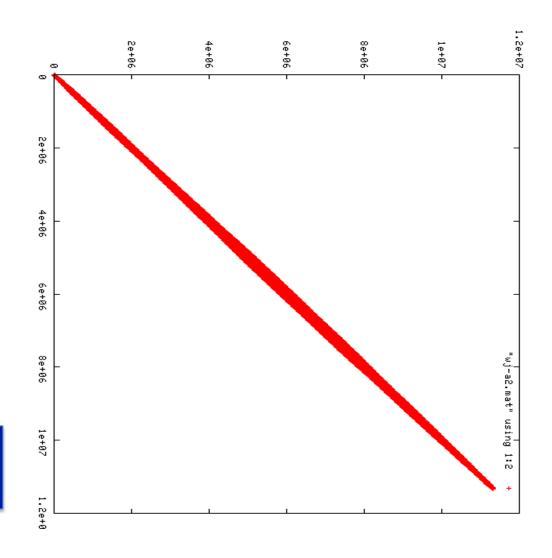
#### System size:

N = 11,333,520

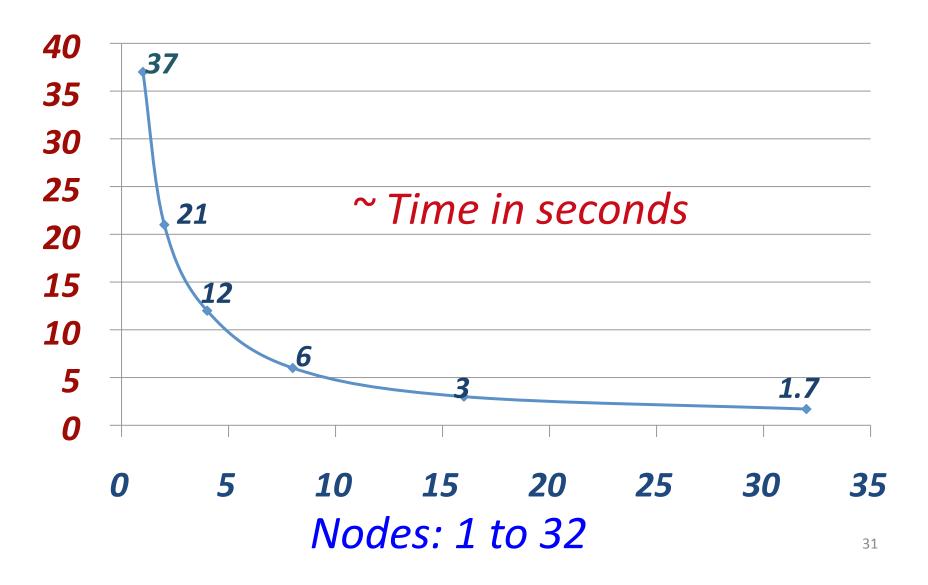
# of nonzeros: 61,026,416

#### <u>bandwidth:</u> 334,613

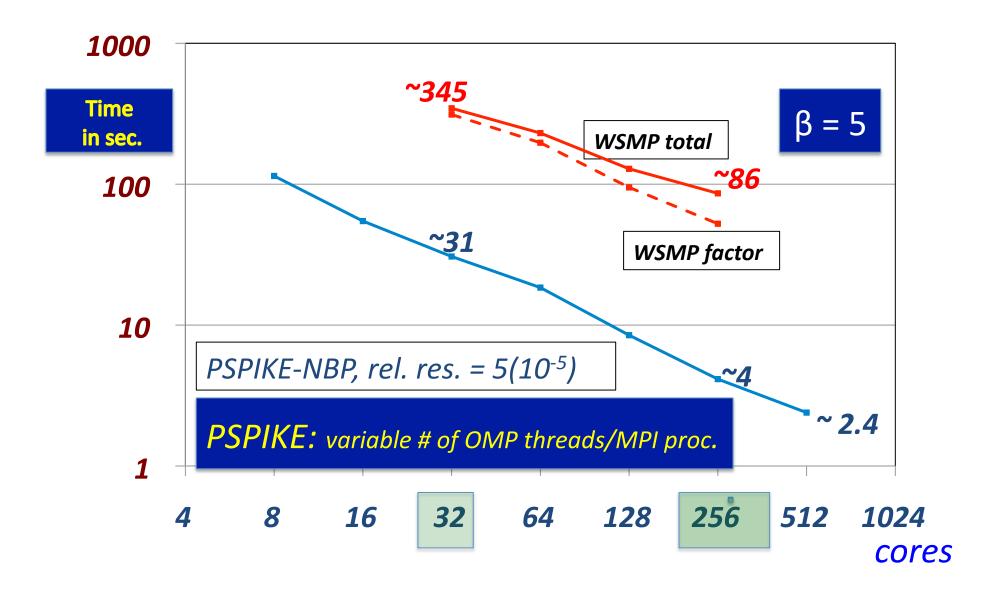
stopping criterion: rel. res. = O(10<sup>-2</sup>)



#### Scalability of TraceMin-Fiedler



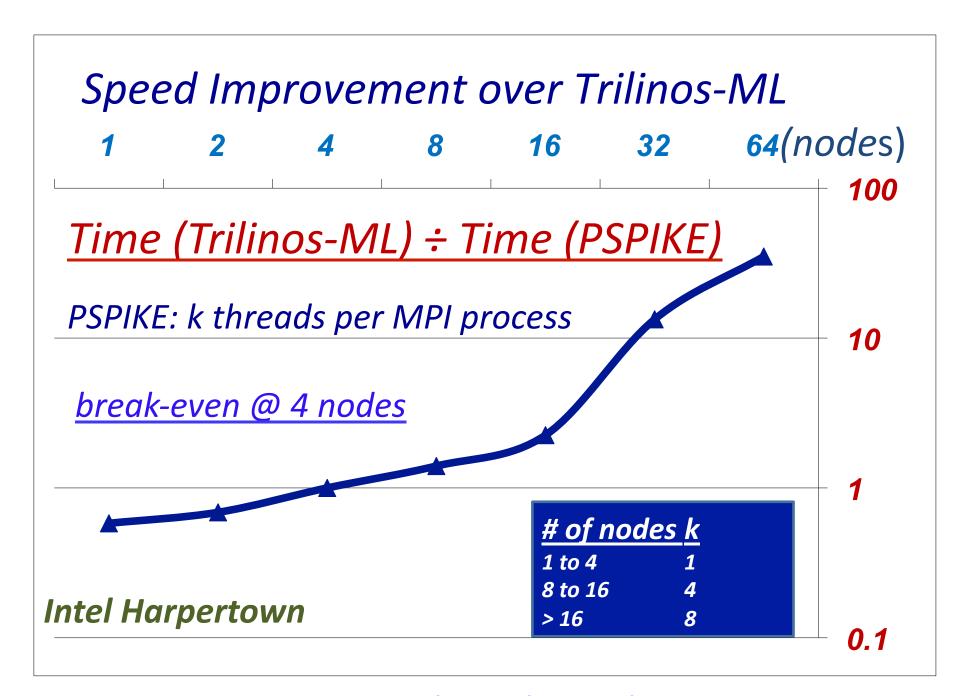
#### MEMS Simulation – benchmark-1



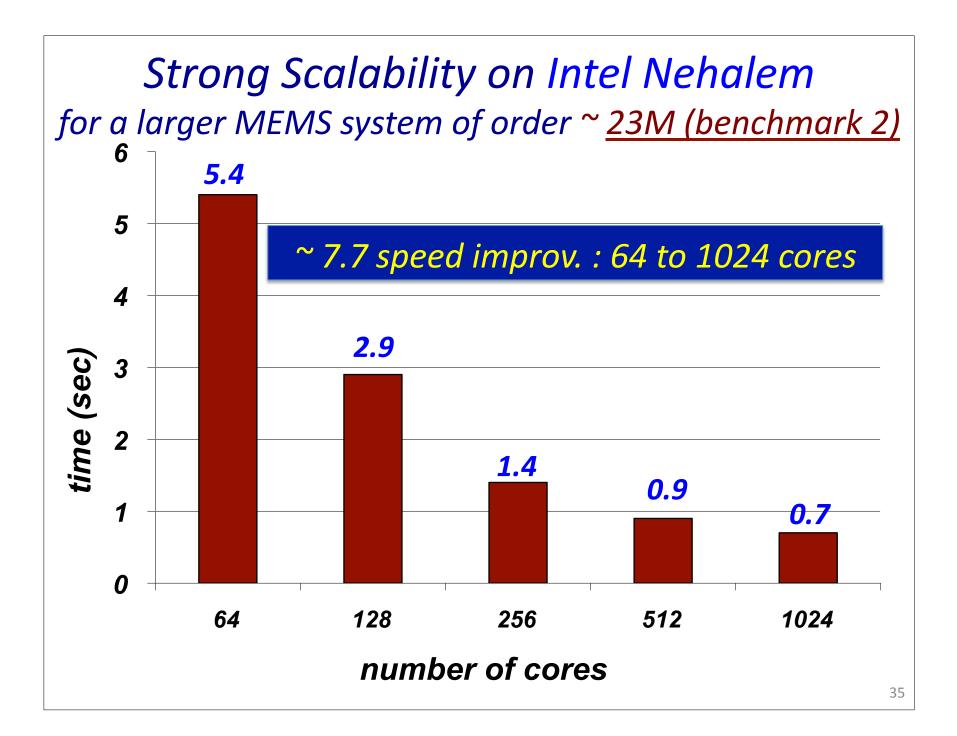
## Scalability of PSPIKE vs. Trilinos Intel Harpertown

- Strong scalability of PSPIKE
   Fixed problem size 1 to 64 nodes (or 8 to 512 cores)
- Comparison with AMG-preconditioned Krylov subspace solvers in:
  - Hypre (LLNL)
  - Trilinos-ML (Sandia)
    - Smoother
      - Chebyshev
      - Jacobi
      - Gauss-Seidel

fastest



MEMS benchmark 1

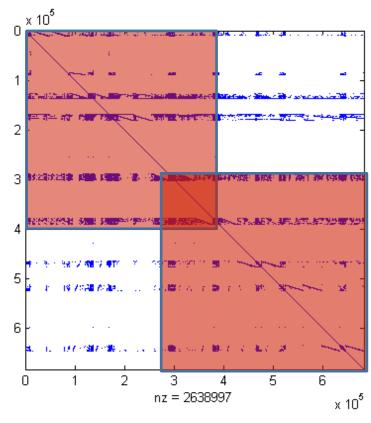


## b. PSPIKE-WBP

#### PSPIKE for ASIC680K system

#### relative residual

Pardiso:  $O(10^{-9})$ ILU+BiCGstab:  $O(10^{-7})$ PSPIKE-WBP:  $O(10^{-7})$ 



Preconditioner consisting of two overlapped blocks (overlap = 11)

Time on 1 node:

Pardiso

– 23 (sec).

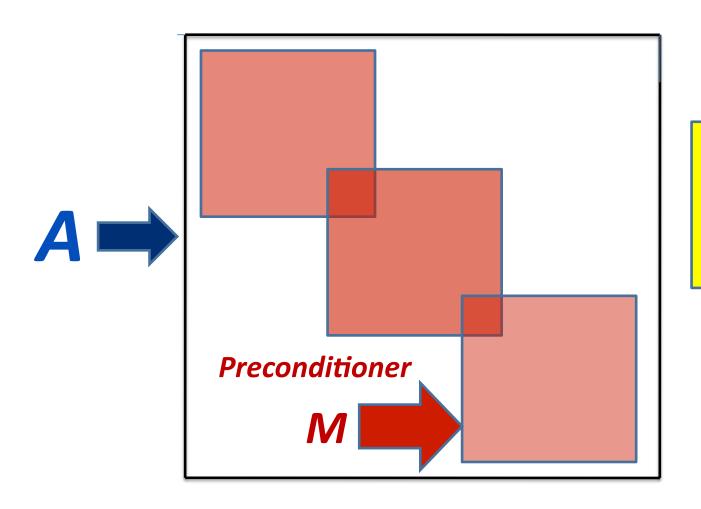
*ILU-BiCGstab* 

-34 (sec).

Time on 2 nodes:

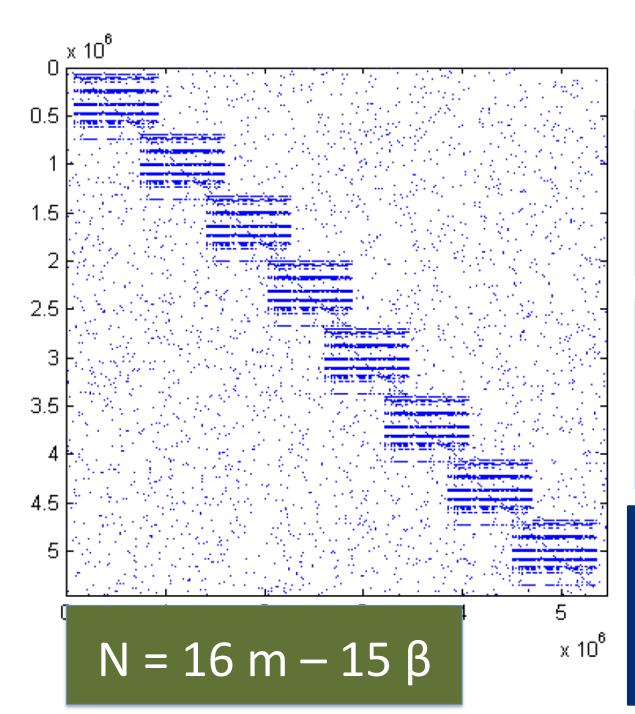
**PSPIKE** 

- 6 (sec).



Overlapped diagonal blocks

PSPIKE-WBP



#### PSPIKE-WBP

ASIC\_680k\_8 nonsymmetric size: ~ 5.5 M nnz: ~31.0 M

Preconditioner: 16 overlapped diagonal blocks overlap size = 100

# PSPIKE-WBP vs. ILU-preconditioned BiCGstab

- ILUT-BiCGstab (one core):
  - -drop tolerance = 10<sup>-3</sup>
  - fill\_in per row = 10%
  - -16 iterations, rel. res. =  $O(10^{-6})$
  - -Time ~ 276 sec.
- PSPIKE-WBP :
  - -1 iteration, rel. res =  $O(10^{-9})$

## PSPIKE-WBP (multiple nodes) vs.

#### ILU-preconditioned BiCGstab (one core)

| # of cores                                 | 32<br>(4 nodes) | 64<br>(8 nodes) | 128<br>(16 nodes) |
|--|-----------------|-----------------|-------------------|
| Time (sec.)                                | 113             | 31              | 9                 |
| Speed improv. over ILU-BiCGstab (one core) | 2.4             | 8.9             | 30.7              |

## Thank You!