# An optimal Q-OR Krylov subspace method for solving linear systems 

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Many Krylov methods have been proposed over the years for solving linear systems $A x=b$

Many of them can be classified as quasi-orthogonal (Q-OR) or quasi-minimum residual (Q-MR)

Q-OR: FOM, BiCG, Hessenberg, ...
Q-MR: GMRES, QMR, CMRH, ...

Whatever their definition, these methods share many fundamental properties

See M. Eiermann and O.G. Ernst, Geometric aspects in the theory of Krylov subspace methods, Acta Numerica, v 10 n 10 (2001), pp. 251-312

They differ by the basis of the Krylov space that is constructed:

- orthogonal for FOM/GMRES,
- bi-orthogonal for BiCG/QMR,
- based on an LU factorization for Hessenberg/CMRH

We use $x_{0}=0$ and assume that $\|b\|=1$, the Krylov space is

$$
\mathcal{K}=\left\{b, A b, \ldots, A^{n-1} b\right\}
$$

## Q-OR methods

We assume that the vectors spanning $\mathcal{K}$ are linearly independent and that we have a basis $V$ of the Krylov space (with columns of unit norm) such that $K=V U$ with

$$
K=\left(\begin{array}{lllll}
b & A b & A^{2} b & \cdots & A^{n-1} b
\end{array}\right)
$$

$V$ nonsingular with $v_{1}=b$ and $U$ upper triangular
We define $H=U C U^{-1}$, upper Hessenberg, where $C$ is the companion matrix for the eigenvalues of $A$. We have $A K=K C$. As a consequence $A V=V H$. It yields an Arnoldi-like relation

$$
A V_{k}=V_{k} H_{k}+h_{k+1, k} v_{k+1} e_{k}^{T}
$$

where $V_{k}$ is the matrix of the $k$ first columns of $V$ and $H_{k}$ is the $k \times k$ principal matrix of $H$

The iterates are defined as

$$
x_{k}=V_{k} y^{(k)}
$$

The residual $r_{k}=b-A x_{k}$ is

$$
\begin{aligned}
r_{k} & =V_{k} e_{1}-A V_{k} y^{(k)} \\
& =V_{k}\left(e_{1}-H_{k} y^{(k)}\right)-h_{k+1, k} y_{k}^{(k)} v_{k+1}
\end{aligned}
$$

The Q-OR method is defined (provided that $H_{k}$ is nonsingular) by

$$
H_{k} y^{(k)}=e_{1}
$$

This annihilates the first term in the residual and the residual norm is $h_{k+1, k}\left|y_{k}^{(k)}\right|$

## Properties of Q-OR methods

Let $r_{k}^{O}$ be the residual vectors of the Q-OR method
Whatever the basis is we can show by induction that

$$
\left|\left(U^{-1}\right)_{1, k}\right|=\left|\nu_{1, k}\right|=\frac{1}{\left\|r_{k-1}^{O}\right\|}
$$

The inverses of the Q-OR residual norms can be read from the first row of the inverse of $U$ (remember that $K=V U$ )

For this property and more see:
G. Meurant and J. Duintjer Tebbens, On the convergence of $Q-O R$ and Q-MR Krylov methods for solving nonsymmetric linear systems, BIT Numerical Mathematics, v 56 n 1 (2016), pp. 77-97

## Construction of "good" bases

We would like to find bases which lead to a "good" (or even optimal) convergence of the Q-OR method

- The matrix $V$ of the basis is related to the Krylov matrix $K$ by $K=V U$ with $U$ upper triangular
- The entries of the first row of $U^{-1}$ are the inverses of the Q-OR residual norms (up to the sign)

Constructing a "good" basis may seem easy since one can think that we can just construct any upper triangular matrix $U^{-1}$ with entries of large modulus on the first row
But, it is not so since the columns of $V$ have to be of unit norm
Moreover, it is not recommended to use the matrix $U$ numerically

## A non-orthogonal optimal basis

Can we construct a basis such that Q-OR minimizes the residual norms?

We would like to construct $H$ column by column without using $U$. We have

$$
H_{j}=U_{j} E_{j} U_{j}^{-1}+\left(\begin{array}{llll}
0 & \cdots & 0 & \frac{1}{u_{j, j}} U_{1: j, j+1}
\end{array}\right)
$$

$E_{j}$ down-shift matrix
It yields

$$
\sum_{j=1}^{k+1} \nu_{1, j} h_{j, k}=0 \Rightarrow \nu_{1, k+1}=-\frac{1}{h_{k+1, k}} \sum_{j=1}^{k} \nu_{1, j} h_{j, k}
$$

At step $k$ we have already computed $\nu_{1, j}, j=1, \ldots, k$ and we would like to choose $h_{j, k}, j=1, \ldots, k+1$ to maximize the absolute value of $\nu_{1, k+1}$ But $h_{k+1, k}$ has to be chosen to obtain a vector $v_{k+1}$ of unit norm Let

$$
\tilde{v}=A v_{k}-\sum_{j=1}^{k} h_{j, k} v_{j}
$$

the next basis vector is $v_{k+1}=\tilde{v} / h_{k+1, k}$ with $h_{k+1, k}=\|\tilde{v}\|$

$$
\left|\nu_{1, k+1}\right|=\frac{\left|\nu^{\top} y\right|}{\|d-B y\|}
$$

with

$$
\begin{gathered}
d=A v_{k}, \quad B=V_{k}=\left(\begin{array}{lll}
v_{1} & \cdots & v_{k}
\end{array}\right), \quad y=\left(\begin{array}{lll}
h_{1, k} & \cdots & h_{k, k}
\end{array}\right)^{T} \\
\nu=\left(\begin{array}{lll}
\nu_{1,1} & \cdots & \nu_{1, k}
\end{array}\right)
\end{gathered}
$$

We need to minimize $1 /\left|\nu_{1, k+1}\right|^{2}$

We would like to solve

$$
\gamma_{o p t}=\min _{y \in \mathbb{R}^{k}, \nu^{T}} \frac{\|d-B y\|^{2}}{\left(\nu^{T} y\right)^{2}}
$$

The minimum is given by

$$
\gamma_{\mathrm{opt}}=\frac{\alpha}{\alpha \nu^{T}\left(B^{T} B\right)^{-1} \nu+\omega^{2}}
$$

with $\alpha=d^{T} d-d^{T} B\left(B^{T} B\right)^{-1} B^{T} d$ and $\omega=d^{T} B\left(B^{T} B\right)^{-1} \nu$
Moreover, if $\omega \neq 0$, a solution $y_{\text {opt }}$ of the minimization problem is given by

$$
\begin{aligned}
y_{\text {opt }} & =\left(B^{T} B\right)^{-1} B^{T} d+\frac{\alpha}{\omega}\left(B^{T} B\right)^{-1} \nu \\
& =s+\frac{\alpha}{\omega} t
\end{aligned}
$$

In our case for computing the solution we have to solve

$$
\left(V_{k}^{T} V_{k}\right) s=V_{k}^{T} A v_{k}, \quad\left(V_{k}^{T} V_{k}\right) t=\nu
$$

## Properties of the optimal basis

$$
\begin{gathered}
V_{k+1}^{T} v_{k+1}=\frac{1}{\nu_{1, k+1}}\left(\begin{array}{c}
\nu_{1,1} \\
\vdots \\
\nu_{1, k} \\
\nu_{1, k+1}
\end{array}\right) \\
V_{k}^{T} V_{k}=\left(\begin{array}{cccc}
1 & \frac{1}{\nu_{1,2}} & \frac{1}{\nu_{1,3}} & \cdots \\
\frac{1}{\nu_{1, k}} \\
\frac{1}{\nu_{1,2}} & 1 & \frac{\nu_{1,2}}{\nu_{1,3}} & \cdots \\
\frac{l_{1,2}}{\nu_{1, k}} \\
\frac{\nu_{1,3}}{\nu_{1,3}} & \frac{\nu_{1,2}}{\nu_{1,3}} & 1 & \cdots \\
\vdots & \vdots & & \frac{\nu_{1,3}}{\nu_{1, k}} \\
\frac{1}{\nu_{1, k}} & \frac{\nu_{1,2}}{\nu_{1, k}} & \cdots & \\
\vdots & & 1
\end{array}\right)
\end{gathered}
$$

When the method converges, the basis is more and more orthogonal

The inverse of $V_{k}^{T} V_{k}$ is tridiagonal and the matrix $V_{k}^{T} A V_{k}$ is upper triangular

It means that we have constructed a right-conjugate direction method

Moreover

$$
t=\left(V_{k}^{T} V_{k}\right)^{-1} \nu=\left(\begin{array}{c}
0 \\
\vdots \\
0 \\
\nu_{1, k}
\end{array}\right)
$$

This relation is used to simplify the construction of the basis vectors

We can simplify the formulas for the new vector

$$
\omega=t^{T} V_{k}^{T} A v_{k}=\nu_{1, k} v_{k}^{T} A v_{k}
$$

Let $y_{\text {opt }}=s+\frac{\alpha}{\omega} t$

$$
\begin{aligned}
\tilde{v} & =A v_{k}-V_{k} y_{o p t} \\
& =A v_{k}-V_{k} s-\frac{\alpha}{\omega} V_{k} t \\
& =A v_{k}-V_{k} s-\frac{\alpha}{\omega} \nu_{1, k} v_{k} \\
& =A v_{k}-V_{k} s-\frac{\alpha}{v_{k}^{T} A v_{k}} v_{k}
\end{aligned}
$$

and

$$
h_{1: k, k}=s+\beta e_{k}, \quad \beta=\frac{\alpha}{v_{k}^{T} A v_{k}}
$$

## The Q-OR optimal algorithm

We compute incrementally the inverses of the Cholesky factors of $V_{k}^{T} V_{k}$

Let $v_{k}^{A}=A v_{k}$
1- $v_{k}^{V}=V_{k-1}^{T} v_{k}, \quad v_{k}^{t A}=V_{k}^{T} v_{k}^{A}$
2- $\ell_{k}=\tilde{L}_{k-1} v_{k}^{V}, y_{k}^{T}=\ell_{k}^{T} \tilde{L}_{k-1}$
3- if $\ell_{k}^{T} \ell_{k}<1, \ell_{k, k}=\sqrt{1-\ell_{k}^{T} \ell_{k}}$, else $\left(p_{k}^{v}\right)^{T}=y_{k}^{T} V_{k-1}^{T}$,
$\ell_{k, k}=\left\|v_{k}-p_{k}^{v}\right\|$ end

4-

$$
\tilde{L}_{k}=\left(\begin{array}{cc}
\tilde{L}_{k-1} & 0 \\
-\frac{1}{\ell_{k, k}} y_{k}^{T} & \frac{1}{\ell_{k, k}}
\end{array}\right)
$$

$5-\ell_{A}=\tilde{L}_{k} v_{k}^{t A}, \quad s=\tilde{L}_{k}^{T} \ell_{A}$
6- $\alpha=\left(v_{k}^{A}\right)^{T} v_{k}^{A}-\ell_{A}^{T} \ell_{A}, \quad \beta=\frac{\alpha}{\left(v_{k}^{ \pm A}\right)_{k}}$
7-

$$
h_{1: k, k}=\left(\begin{array}{c}
h_{1, k} \\
\vdots \\
h_{k, k}
\end{array}\right)=s+\beta e_{k}
$$

8-

$$
\begin{gathered}
\tilde{v}=v_{k}^{A}-V_{k} h_{1: k, k}, h_{k+1, k}=\|\tilde{v}\|, \nu_{1, k+1}=-\frac{1}{h_{k+1, k}} \nu^{T} h_{1: k, k} \\
\nu=\left(\begin{array}{lll}
\nu_{1,1} & \cdots & \nu_{1, k+1}
\end{array}\right)^{T}
\end{gathered}
$$

9- $v_{k+1}=\frac{1}{h_{k+1, k}} \tilde{v}$ and $v_{k+1}^{A}=A v_{k+1}$
10- if needed, solve $H_{k} y^{(k)}=\|b\| e_{1}$ using Givens rotations, $x_{k}=V_{k} y^{(k)}$

In this algorithm almost everything is expressed in terms of matrix-vector products

## Numerical experiments

SUPG scheme (Streamwise upwind Galerkin)
Convection-diffusion equation in a square with a mesh size of $1 / 41$ The diffusion coefficient is $\nu=0.01$
This matrix is of order 1600 and has 13924 non zero entries. Its norm is $4.871610^{-2}$ and the condition number is 40.518


Difference of the true residual norms of GMRES-MGS and Q-OR optimal, supg 1600


True residual norms of GMRES-MGS (blue) and Q-OR-opt (red), supg 1600, $n=1600$

True residual norms for $k=200$

- GMRES-CGS $1.5404310^{-13}$
- GMRES-CGS with reorthogonalization $7.0558510^{-15}$
- GMRES-CGS with double reorthogonalization $7.2379010^{-15}$
- GMRES-MGS $1.3377610^{-14}$
- GMRES-MGS with reorthogonalization $6.7064910^{-15}$
- GMRES-MGS with double reorthogonalization $6.7033910^{-15}$
- GMRES-Householder 2.03961 10-14
- QOR opt $5.5062610^{-15}$


## Could we use the fact that $\left(V_{k}^{\top} V_{k}\right)^{-1}$ is tridiagonal?

We can compute the non-zero entries of $\left(V_{k}^{\top} V_{k}\right)^{-1}$
Previously we used the Cholesky factors of the matrix $\left(V_{k}^{T} V_{k}\right)^{-1}$ to solve $V_{k}^{T} V_{k} s=v_{k}^{t A}$

In theory these factors are bidiagonal matrices. However, we computed all their entries for the sake of numerical stability

Now we would like to investigate to what extent we can use the fact that $\left(V_{k}^{T} V_{k}\right)^{-1}$ is tridiagonal to compute the vector $s$

This would save many dot products

fs 680 1c, true residual norms, GMRES -MGS(plain) and Q-OR-opt-trid (dashed)

The problem is that the values of $\nu_{1, j}$ which are used to compute the inverse of the matrix $V_{k}^{\top} V_{k}$ are not directly linked to the computed vectors $v_{j}$
After a while there is a discrepancy between $\left(V_{k}^{\top} V_{k}\right)^{-1}$ and what is computed with the $\nu_{1, j}$ 's
We can compute the relative residual norms in two ways:
The first one is $1 /\left|\nu_{1, k+1}\right|$
Let $\tilde{r}_{0}$ be the residual at the beginning of the cycle, the second way of computing the relative residual norm is obtained from solving $H_{k} y^{(k)}=\|b\| e_{1}$ which gives $h_{k+1, k}\left|y_{k}^{(k)}\right| /\left\|\tilde{r}_{0}\right\|$
A simple remedy to our problems is to restart the algorithm when there is a too large difference between the two ways of estimating the relative residual norms

supg 1600, true residual norms, GMRES -MGS(plain) and Q-OR-opt-trid with restart $\varepsilon_{\nu}=10^{-2}$ (dashed)

However, things are not always so nice...

fs 680 1c, true residual norms, GMRES -MGS(plain) and Q-OR-opt-trid with restart $\varepsilon_{\nu}=10^{-2}$ (dashed)

## Conclusion

Using the properties of the Q-OR methods we were able to construct a non-orthogonal basis for which Q-OR gives the same residual norms as GMRES

The algorithm is slightly more expensive than GMRES but it can be simplified using automatic restarts

It is more parallel than GMRES-MGS and most of the operations are matrix-vector products

In many cases the maximum attainable accuracy is better than with GMRES-MGS

It remains to study its stability in finite precision arithmetic and to see how to use it on parallel computers

